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HP-UX 11.X  
Technology and Beyond

***DAL***

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# HP-UX 11.X Technology and Beyond Program Outline

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- Introductions and Overview
- "Largeness" Extensions to HP-UX :
  - Filesystems, Files, UIDs, File Descriptors, Process Space
- 64-bit HP-UX
- Developing for 64-bits
- Kernel Threads
- Additional technologies:
  - DLKM, Large Pages, Ignite-UX, System Recovery, others
- Networking enhancements
- HP-UX 11.X and Extension Pack releases
- Review

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# HP-UX 11.X Technology Review

## Program Outline

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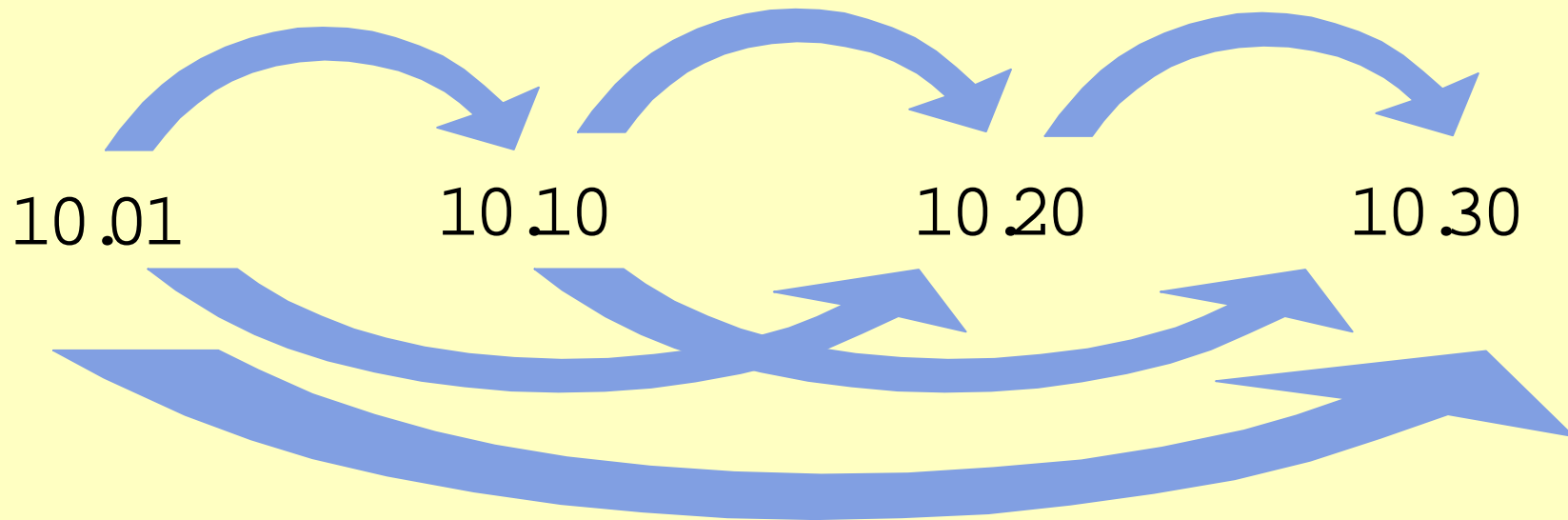
- Technical level of presentation
  - familiarity with System Administration tools
  - familiarity with Software Development concepts
  - familiarity with Operating System fundamentals
- Format of presentation
- References for further study
- Questions

# "Largeness" Features in HP-UX :

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## HP-UX 10.X Releases

- Each release is a superset of the previous
- HP-UX 10.01 is the "gateway" to the family
- Upgrade when and if you need



# HP-UX 10.10 : CQ 1 1996

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## Main new features

- 128 GB file system
- 3.75 GB RAM (T500, 0.75 GB cards)
- 1.9 GB process data space
- 60K File descriptors/process
- Shared LVM (SLVM) - for OPS
- Spec 1170 (UNIX 95)
- CDE - Common Desktop Environment
- 4 byte EUC commands
- DHCP server (including SAM management)
- SAM management of NIS
- 48 LAN card support

high-end  
focus

standards  
focus

**SIZE EXTENSIONS**

# HP-UX 10.20 : CQ 3 1996

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## Main new features

- Performance
  - PA-8000 optimization, Fibre Channel net/storage
  - MP tuning of transports, stacks, sockets, drivers
  - LVM tuning
  - Processor Affinity support
- Large files (local) - 128 GB files
- >60K UID s - enable 4 billion user ID s
- 64-bit register math (PA 8000)
- 2.75 GB Shared Memory via patch
- Technical application performance
- DHCP client
- Distributed Print Services (Palladium)

high  
perf

high-end  
focus

standards  
focus

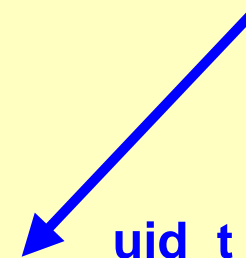
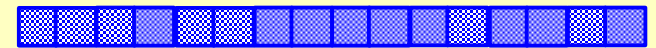
**SIZE EXTENSIONS**

Full performance for PA-8000 based systems

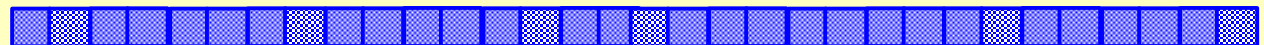
## Large UIDs

- Base type of uid increased from 16 to 32 bits
- Changes to kernel, filesystems, libraries, commands, APIs
  - requires recompile to use Large UIDs
- For large number of users -or- for sparsely mapped uids in a large range (e.g.. telephone numbers)
- HFS: the HP-UX kernel detects and converts HFS filesystems on the fly
- VxFS: supported in Version 3.0 (10.20)

uid\_t (16 bits)



uid\_t (32 bits)



# Using Large UIDs (continued)

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```
< /dev/vg00/lvol4 mounted on /tmp >
% fsadm /dev/vg00/rlvol4
file system      :      /dev/vg00/rlvol4
magic number    :      95014
feature bits    :      1
file system supports : nolargefiles, longfilenames

% touch lg_uid_file
% chown 99999 lg_uid_file
% ls -l lg_uid_file
-rwxrwxrwx  1 99999  sys   0 Jan 21 15:09 lg_uid_file

% fsadm /dev/vg00/rlvol4
file system      :      /dev/vg00/rlvol4
magic number    :      95014
feature bits    :      5
file system supports : nolargefiles, largeuids, longfilenames

<kernel has changed UID structures on the filesystem>
```



# HP-UX 10.30 (limited release)

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## Main new features

- Performance
  - PA-8000 (continued)
- 1x1 kernel threads
- 128 GB files (networked)
- NFS Pv3, NFS+
- libc versioning
- Native OpenGL
- Support for new systems & peripherals
- Stream based TCP/IP
- BIND 4.9.3
- Year 2000 clean
- T600

high  
perf

high-end  
focus

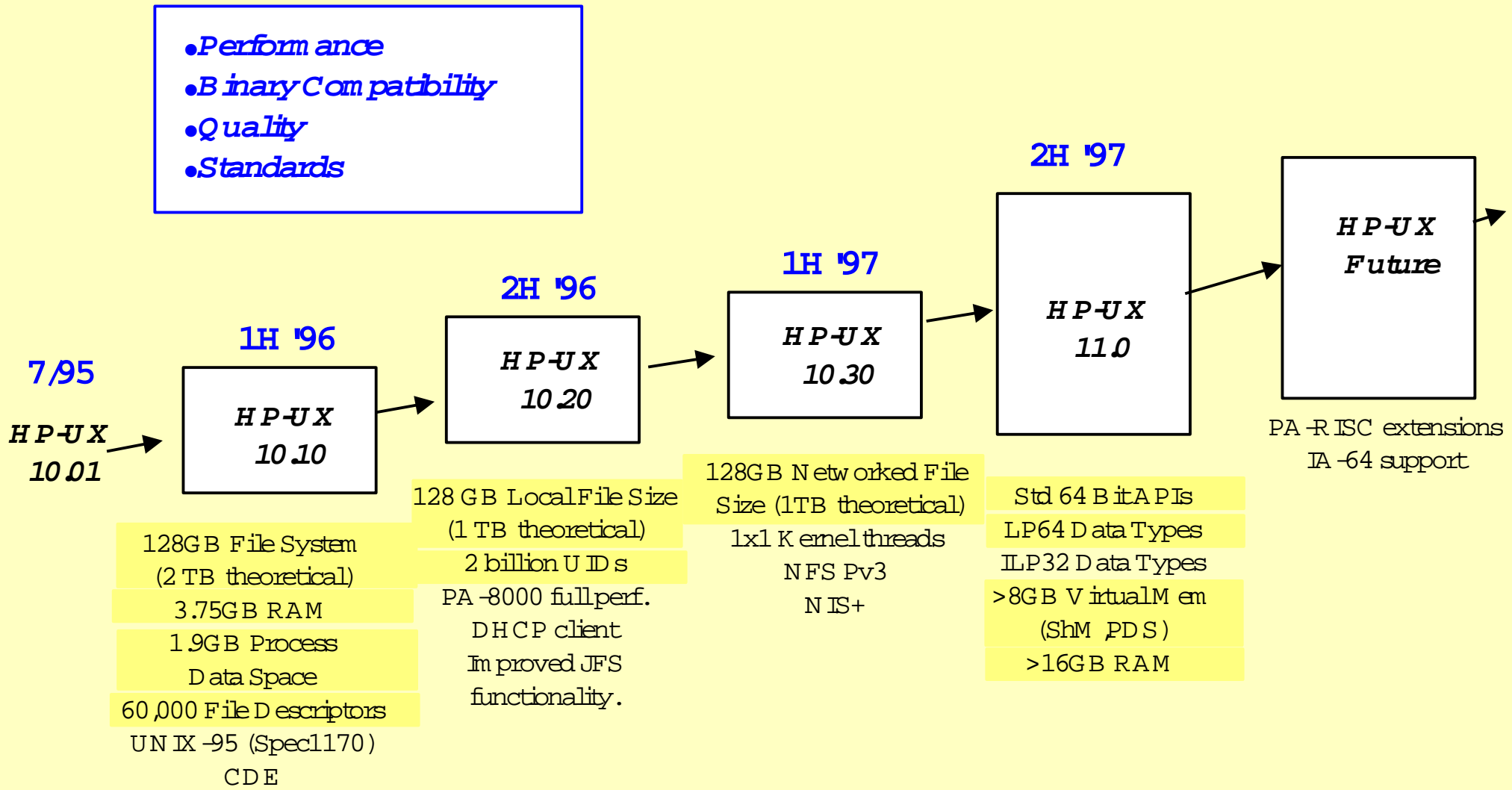
standards  
focus

**SIZE EXTENSIONS**

# HP-UX Operating System Roadmap

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- Performance
- Binary Compatibility
- Quality
- Standards



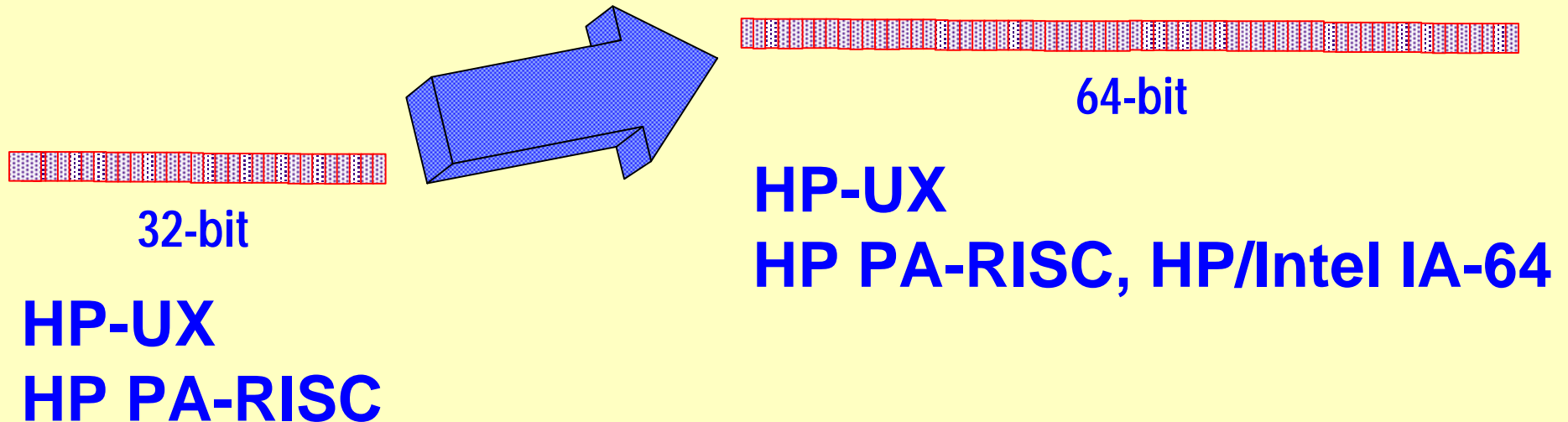
**Extend OS Size Capability**

## Market drivers to 64 bits

- Some Database vendors and technical software developers are leading the movement to 64-bit computing.
- They developed databases and applications that handle very large memory and enable access by many more users.

# HP's 64-bit strategy: overview

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## Key theme:

Implement evolutionary -- not revolutionary -- product strategies to deliver key new features and protect customers' software investments.

# Elements of HP's evolutionary 64-bit strategy

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Objective	Why important	Strategy
Lead industry in hardware and software features and performance	Meets customers' increasing functionality and performance demands	Collaborate with partners in developing next-generation UNIX and chip technologies
Provide smooth upgrade path	Ensures investment protection; minimizes upgrade costs in time and money	Continue policy of forward binary compatibility; 32- and 64-bit apps to coexist and communicate
Continue HP's long-standing commitment to standards	Gives partners flexibility in platform selection	Lead standards development and adherence

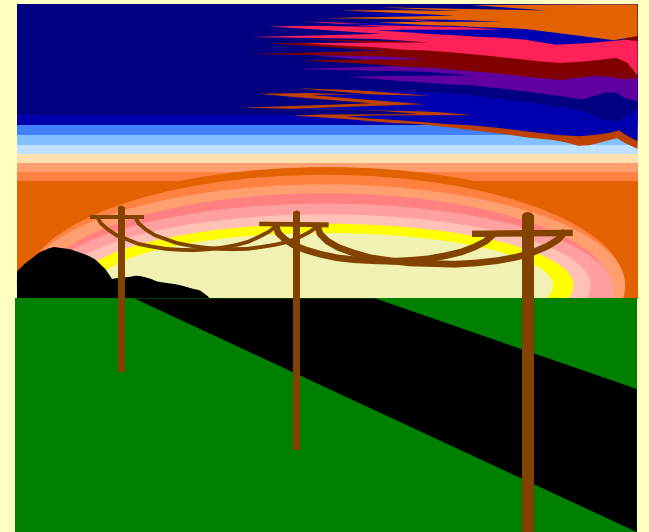
# Benefits of 64-bit computing

- Scalability
  - Larger applications and data and more users
- Potential gains in performance
  - Much larger amounts of data can reside in RAM, resulting in performance gains due to much less (time-consuming) swapping to disk
- These attributes make 64 bits well suited for certain high-end applications
  - very large DB and Decision Support
  - OLTP with 10's of thousands of users
  - complex technical simulations

# Smooth upgrade to 64-bit environment

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- Investment protection through forward binary compatibility
  - 32-bit applications may run unmodified on 64-bit HP-UX
- No migration; minimal end-user effort
  - No forced recompile
  - No forced recode
  - No data reload



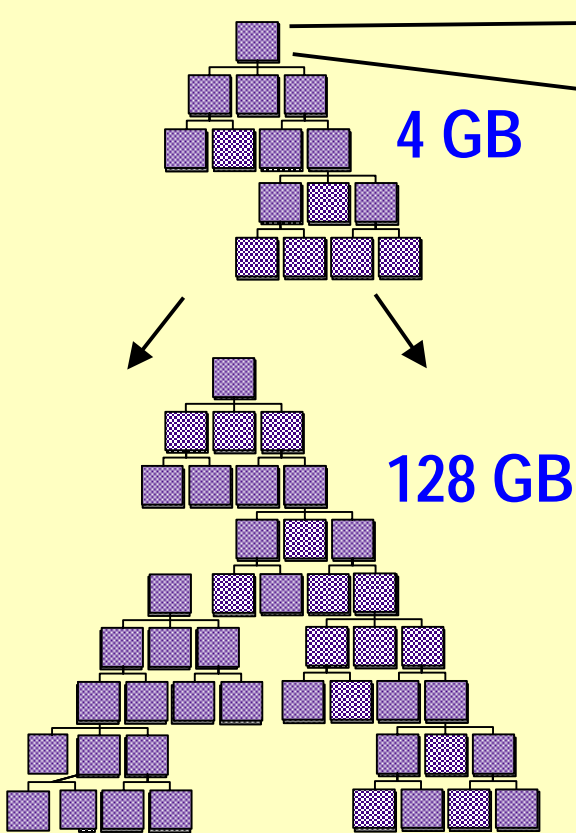
# Evolutionary introduction of 64-bit functionality into HP-UX

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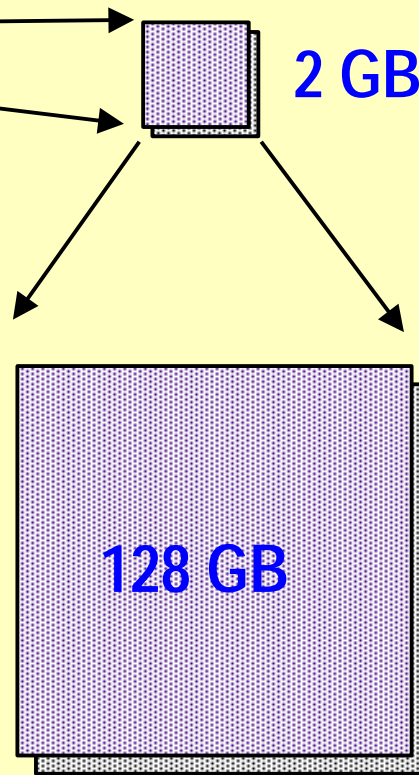
HP-UX 10.10, 2/96

HP-UX 10.20, 8/96

HP-UX 11.00, 1997

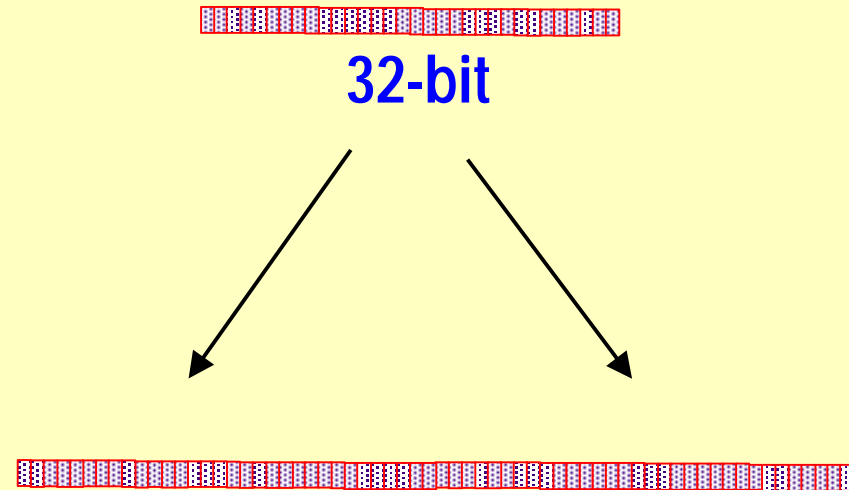


Large file system



Large file size

64-bit register math



Large memory/addressing

64-bit data model



# Operating System Data Models:

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## ILP32 and LP64

	ILP32	LP64
Integer	32 bits	32 bits
Long	32 bits	64 bits
Pointer	32 bits	64 bits

- HP-UX 10.0 is ILP32 as are many other UNIXes
- HP-UX 10.[10,20,30] extended the OS capabilities
- HP-UX 11.0 comes in two versions:
  - ILP32
  - LP64
- can cross-develop between 32- and 64-bit OS versions
- can execute both 32- and 64-bit applications on 64-bit kernel
- can execute only 32-bit applications on 32-bit kernel

# HP-UX Operating System : Specifications by Version

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Attribute	HP-UX 10.01	HP-UX 10.10	HP-UX 10.20	HP-UX 10.30	HP-UX 11.00/32	HP-UX 11.00/64
Introduced	Jun 95	Feb 96	Aug 96	Aug 97	Nov 97	Nov 97
File system	4 GB	128 GB	128 GB	128 GB	128 GB	128 GB
File size	2 GB	2 GB	128 GB local, 2GB network	128 GB local and network	128 GB local and network	128 GB local and network
Physical RAM	2 GB	3.75 GB	3.75 GB	3.75 GB	3.75 GB	4 TB
Shared Mem	1.75 GB	1.75 GB	2.75 GB	2.75 GB	2.75 GB	8 TB
Process data space	0.9 GB	1.9 GB	1.9 GB	1.9 GB	1.9 GB	4 TB
# File Descriptors	2,000	60,000	60,000	60,000	60,000	60,000+
# User IDs	60,000	60,000	2	2	2 billion	2 billion
Threads model	User	User	User	User and Kernel	User and Kernel	User and Kernel
Y2K Ready	patch	patch	patch	yes	yes	yes

# HP-UX 11.0 Specifications

***DAL***

Attribute	32-bit version	64-bit version
CPUs supported	16	32
File system size	128 GB	128 GB
File size (local and networked)	128 GB	128 GB
Physical RAM	3.75 GB	4 TB
Shared memory	2.75 GB	8 TB
Process data space	1.9 GB	4 TB
File descriptors	60,000 plus	60,000 plus
User IDs	2 billion	2 billion
Threads model	User and Kernel	User and Kernel

# HP-UX 11.0 – Supported Systems (Nov 1999 Extension Pack)

**DAL**

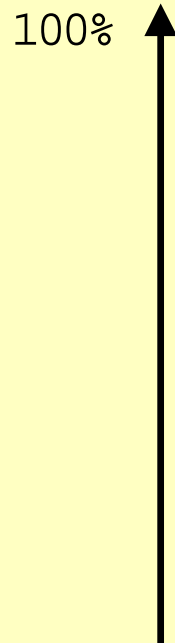
Series	Model	32-bit	64-bit
A-Class	A180, A180C	X	
D-Class	Dx10, Dx20, Dx30, Dx50, Dx60	X	
	Dx70, Dx80, Dx90	X	X
E-Class	E25, E35, E45, E55	X	
F-Class	F10, F20, F30	X	
G-Class	G30, G40, G50, G60, G70	X	
H-Class	H20, H30, H40, H50, H60, H70	X	
I-Class	I30, I40 I50, I60, I70	X	
K-Class	K100, K200, Kx10, Kx20	X	
	Kx50, Kx60, Kx60, Kx70, Kx80	X	X
L-Class	L1000, L2000		X
N-Class	N4000-[36,44]		X
R-Class	R380, R390	X	X
T-Class	T500, T520	X	
	T600	X	X
V-Class	V2200, V2250, V2500		X
EPS	EPS22, EPS23	X	X
	EPS40		X
700 Series	712, 715/[64,80,100,100XC], 725/100	X	
B-Class	B132L, B132L+, B160L, B180L	X	
	B1000		X
C-Class	C100, C110, C160L	X	
	C160, C180[XP], C200, C240, C360	X	X
	C3000		X
J-Class	J200, J210, J210XC	X	
	J280, J282, J2240	X	X
	J5000, J7000		X
J-Class	J280, J282, J2240	X	X

HP-UX users may gain performance increases without needing to recompile their application

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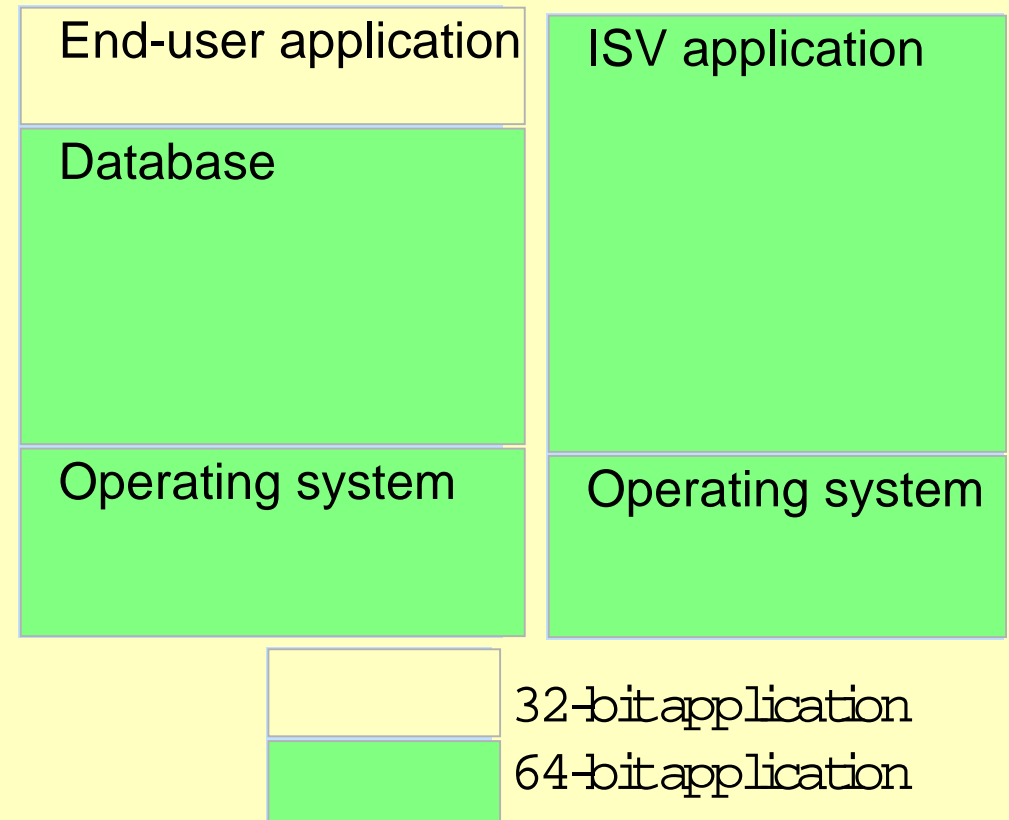
- HP-UX will allow an existing 32-bit end-user application to interact with a 64-bit database
- Most of the total performance gain will come from recompiles of key DB and ISV applications
- Customers need not recompile their applications

Contribution to performance



**Most commercial environments**

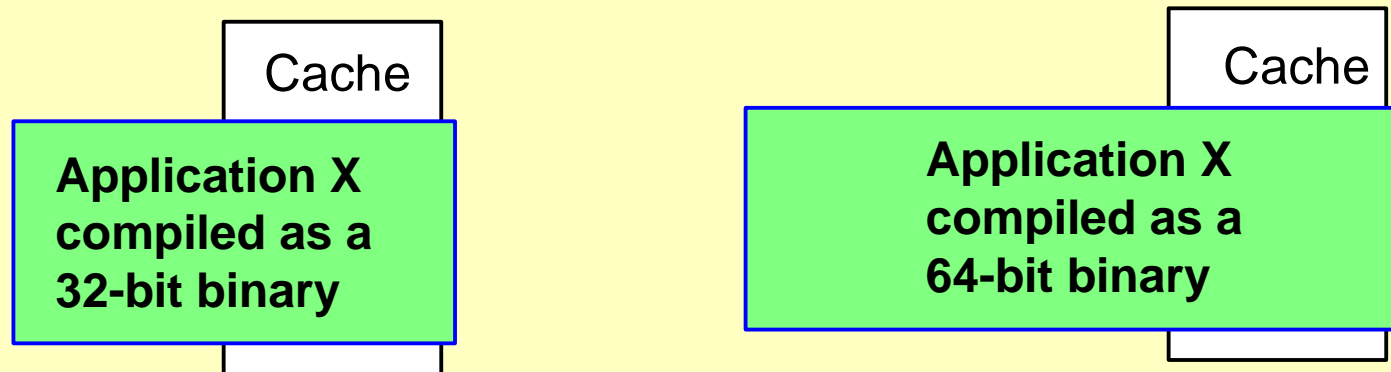
**Many technical environments**



# Mixing 32- and 64-bit applications

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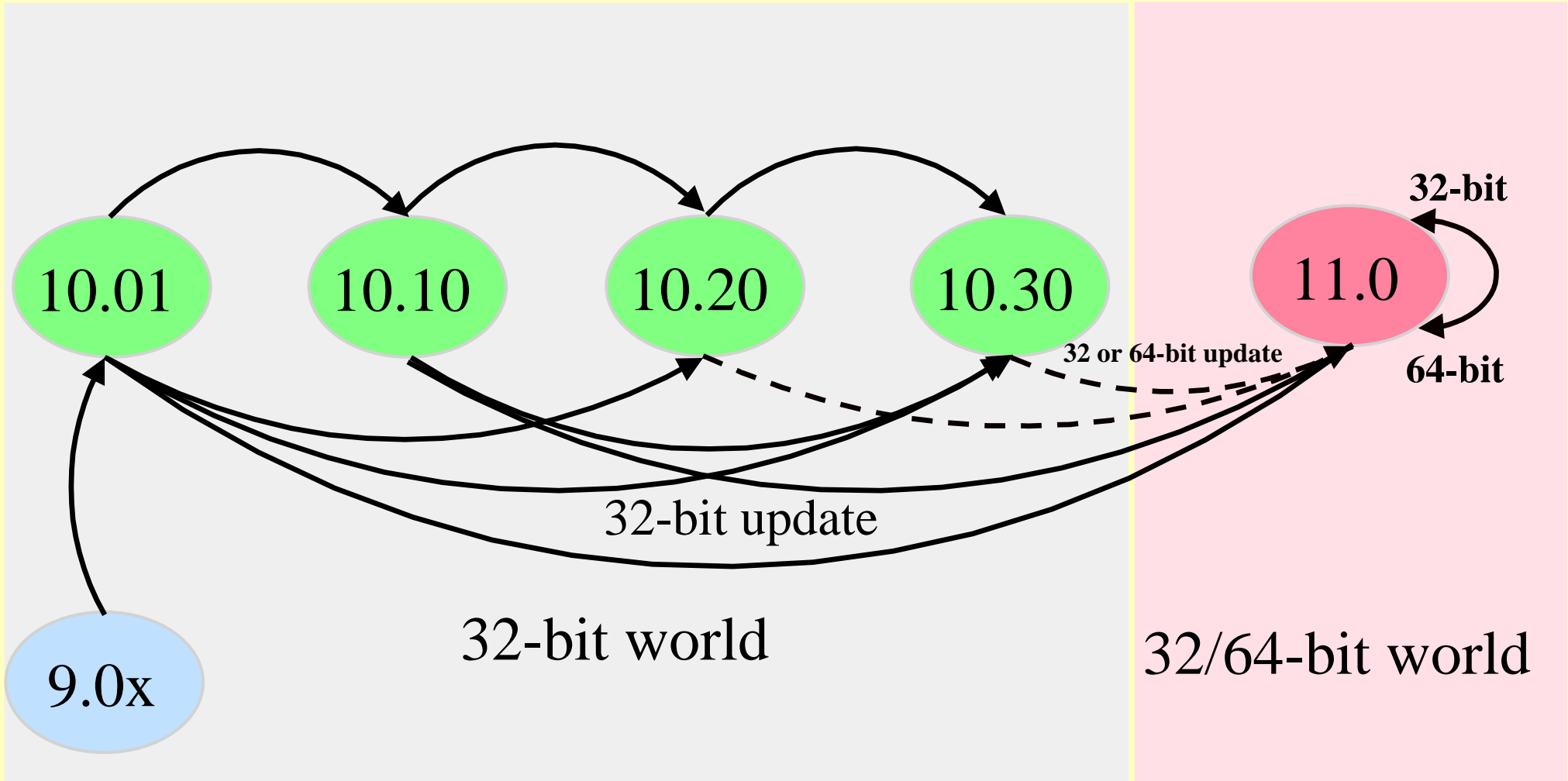
- 32-bit applications may run faster than recompiled 64-bit versions, due to "cache-fit" effect



**Best overall performance comes through 32- and 64-bit coexistence, which HP-UX provides**

# Upgrade Paths to HP-UX 11.0

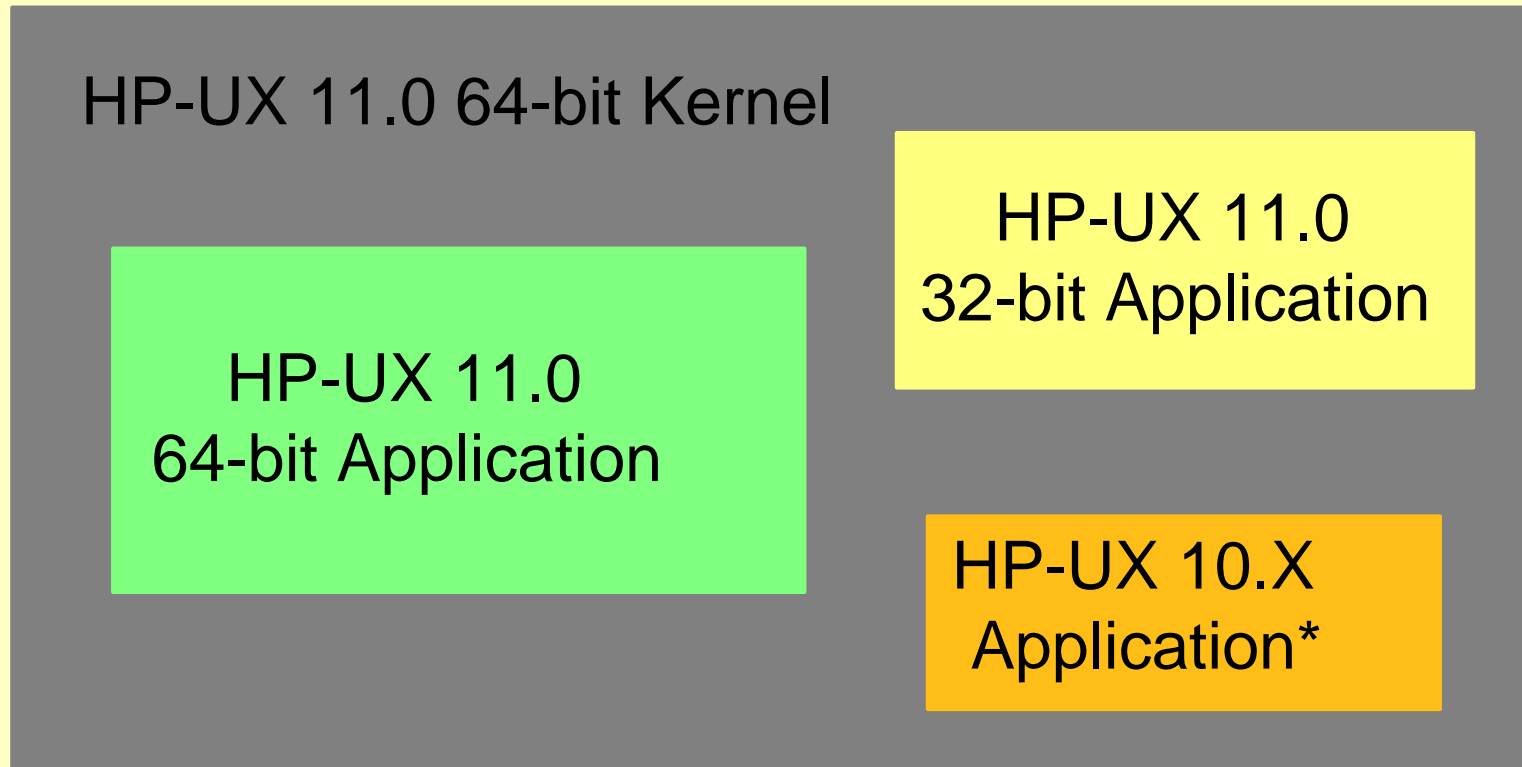
Software Distributor support updates from HP-UX 10.x to HP-UX 11.0 in one step



-- update directly to either 32-bit or 64-bit

# Applications that run on 64-bit HP-UX 11.0

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★ Well-behaved applications that run on HP-UX 10.X run on HP-UX 11.0. (See Compatibility Guidelines)



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# Applications that run on 32-bit HP-UX 11.0

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HP-UX 11.0 32-bit Kernel

HP-UX 11.0  
32-bit Application

HP-UX 10.X  
Application\*

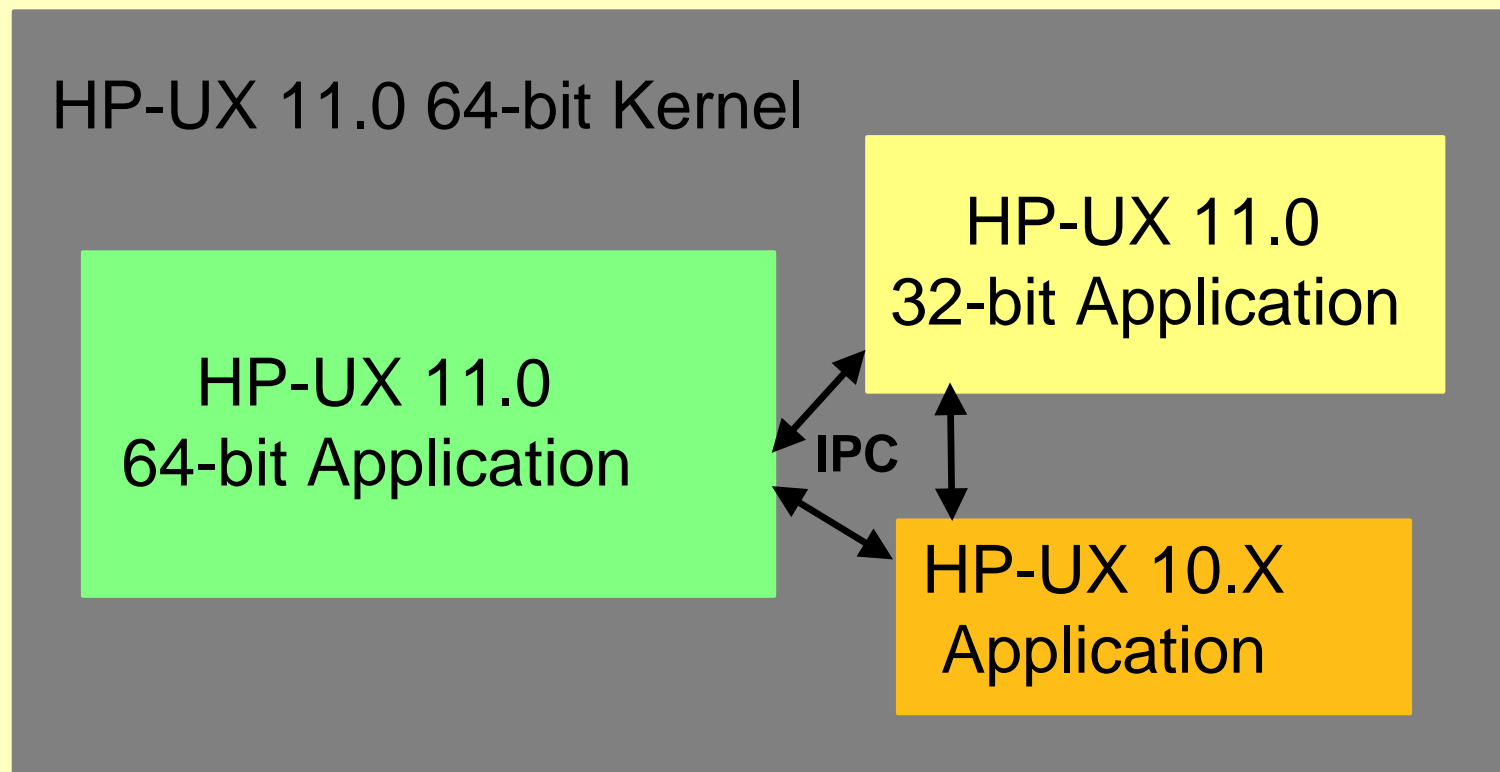
- ★ Well-behaved applications that run on HP-UX 10.X run on HP-UX 11.0. (See Compatibility Guidelines)

# Applications Interoperability on 64-bit HP-UX 11.0

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- 32- and 64-bit applications can interoperate on 64-bit HP-UX using standard IPC mechanisms:

- Shared Memory
- Mapped Files
- Sockets
- Signals
- Message Queues
- Pipes
- RPC



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# HP-UX Compilers

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- Available for both 32- and 64-bit development
  - HP C
  - HP aC++
  - HP FORTRAN 90
  - HP Assembler
  - HP DDE
  - HP PAK
  - HP Linker toolset
- Available only for 32-bit development
  - HP FORTRAN 77
  - HP Pascal
  - HP MicroFocus COBOL
  - HP C++ (cfront)

# Compiler Option/Hardware Architecture Run-time Compatibility

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Compiler Option	PA-RISC 1.1 32-bit Systems	PA-RISC 2.0 64-bit Systems
+DA1.1	X	X
+DAportable	X	X
+DA2.0		X
+DD64 or +DA2.0W		X

+DD64 is a HP C option for compiling in 64-bit mode.

+DA2.0W is the HP aC++, HP Fortran90, and HP C option for compiling in 64-bit mode.

# Compiler Options to Develop for 32- or 64-bit Applications

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Compiler Option	What it Does
+DA1.1	Produce PA1.1 code
+DA2.0	Produce PA2.0 code (32 bit)
+DA2.0W	Produce PA2.0 code (64 bit)
+DD32	same as +DA1.1
+DD64	same as +DA2.0W

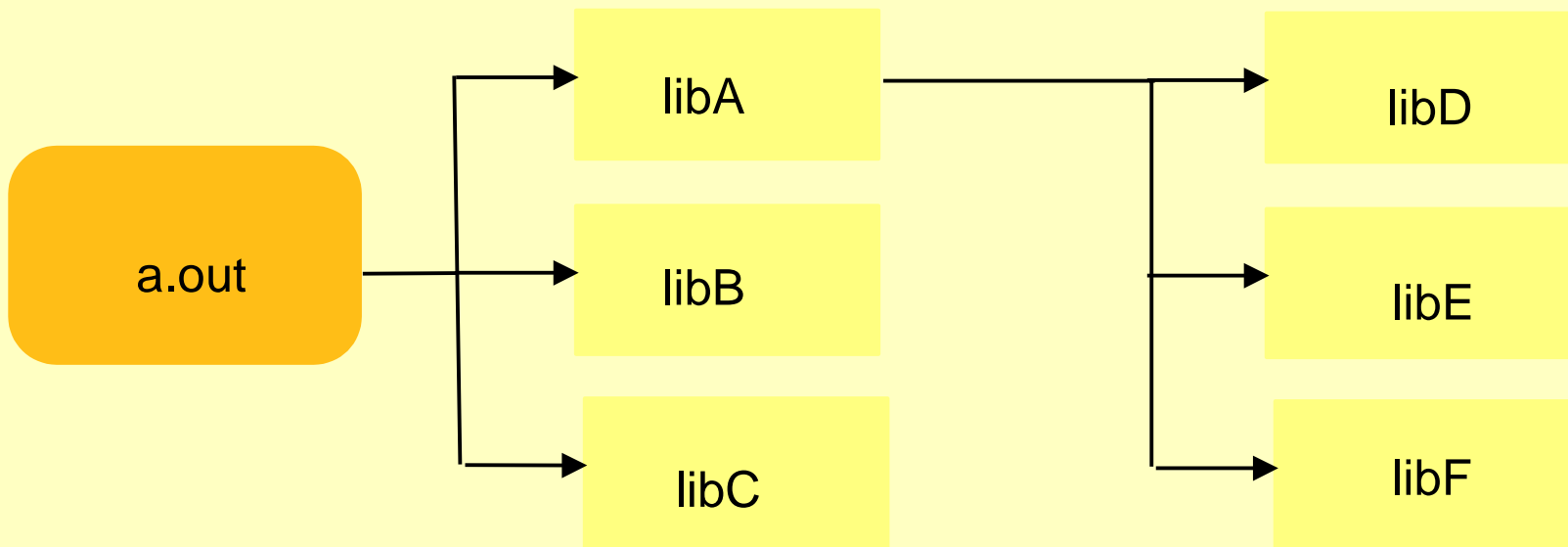
- all options are supported on both 32- and 64-bit systems
  - can cross-develop for either platform
- on 32-bit systems make sure you have installed the 64-bit library fileset (OS-CORE-C-M-IN-64ALIB)

# HP-UX Linker Changes

- A number of new /changed options to HP-UX linker - try "ld +help" or review STK Whitepapers
- Apps compiled and linked in 32-bit mode will see no change in behavior.
- Apps compiled and linked in 64-bit mode use a runtime startup model similar to other SVR4 systems.

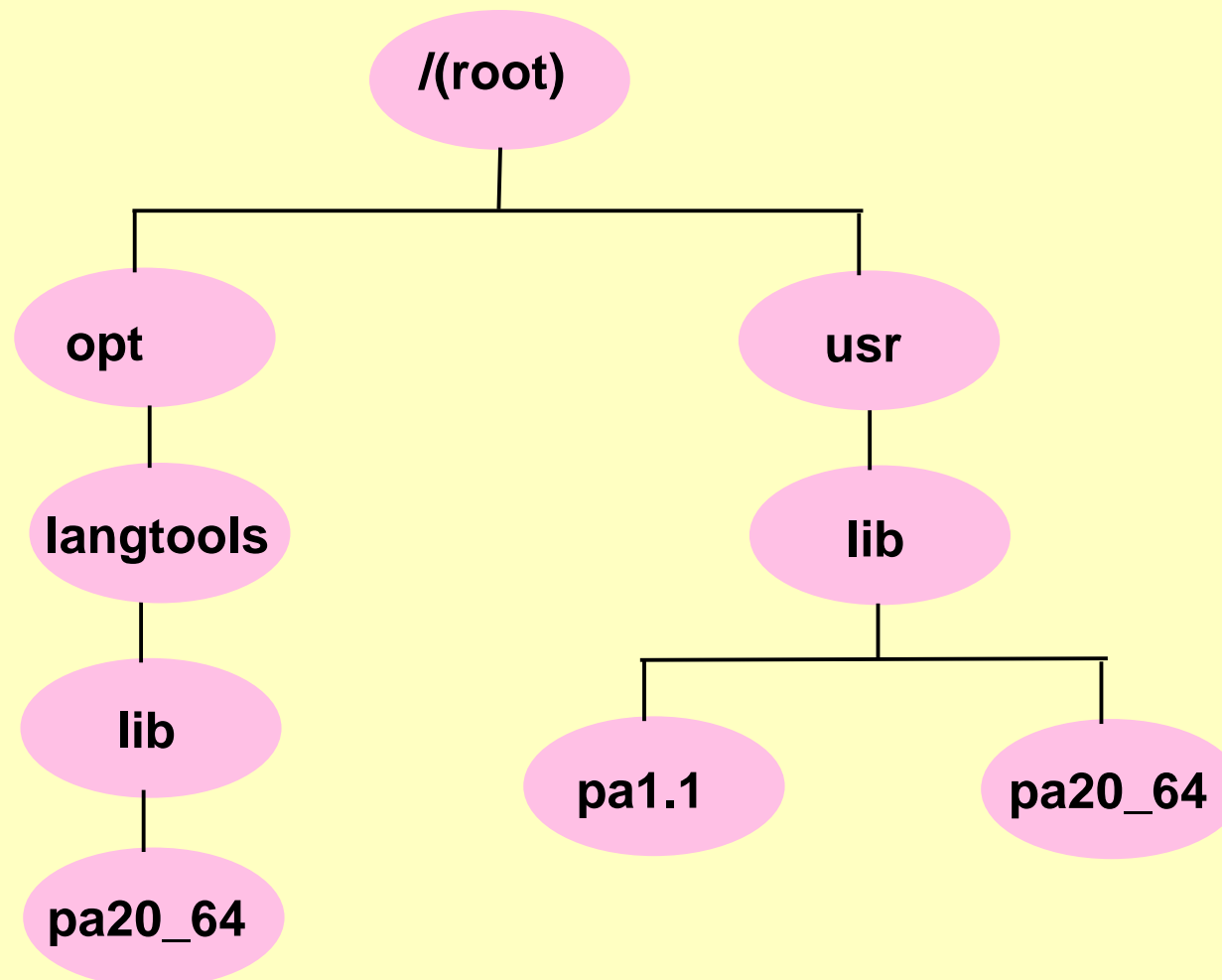
# New Linker search order for 64-bit applications

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- 64-bit mode: Breadth-first search:  
a.out → libA → libB → libC → libD → libE → libF
- 32-bit mode: Depth-first search:  
a.out → libA → libD → libE → libF → libB → libC

# New Library Paths



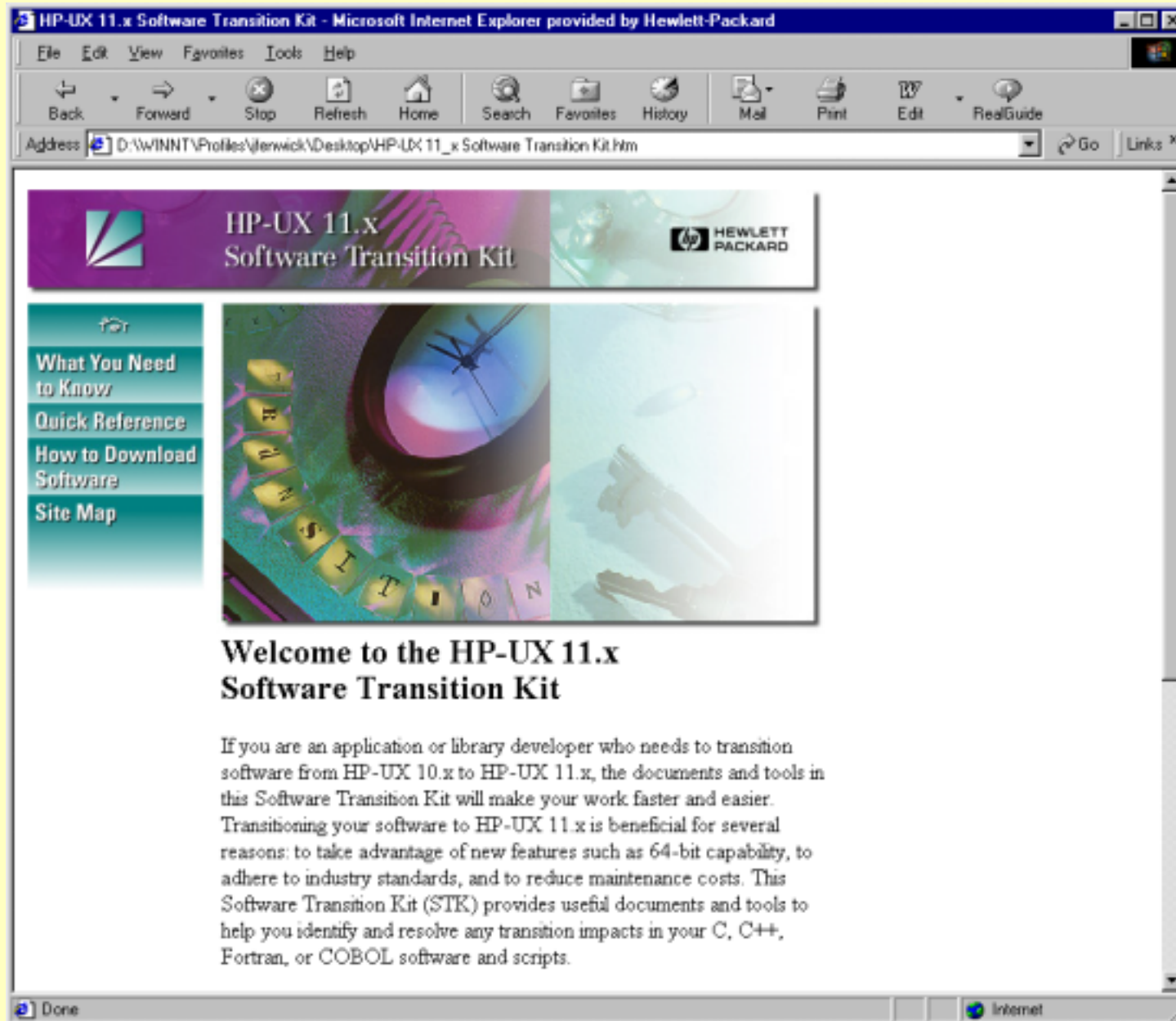


# HP-UX Transition Tools

- HP-UX Software Transition Kit (STK). Current: 1.3.1
  - tools and documentation to discover changed APIs
  - works on C and C++ source code, scripts, makefiles
  - <http://www.software.hp.com/STK/>
- HP C compiler
  - both lint and the C compiler provide options to help transition to the 64-bit data model.
- HP-UX 11.0 Release Notes
  - documents system header file changes, system library changes, and lists 64-bit versions of system libraries.
- FlexeLint
  - identifies non-portable constructs in C and C++ programs.

# HP-UX 11.0 STK

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# HP-UX 32-bit and 64-bit Base Data Types

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Data Type	ILP32 size(bits)	LP64 size(bits)
char	8	8
short	16	16
int	32	32
<b>long</b>	<b>32</b>	<b>64</b>
long long	64	64
<b>pointer</b>	<b>32</b>	<b>64</b>
float	32	32
double	64	64
long double	128	128
enum	32	32

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# ILP32 to LP64

## Porting Concerns

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- Fundamental changes:
  - longs and ints are no longer the same size
  - pointers and ints are no longer the same size
  - pointers and longs are 64 bits and are 64-bit aligned
  - Predefined types `size_t` and `ptrdiff_t` are 64-bit integral types
- Potential impact:
  - data truncation
  - data type promotion
  - constants
  - enumerated types
  - pointers
  - data alignment and data sharing
  - bit shifts and bit masks
  - bit fields

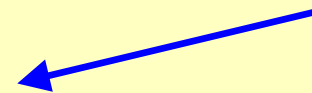
# Example: code works on 32-bit and will fail on 64-bit

***DAL***

```
int main ()
{
    int *buffer;
    buffer = malloc(sizeof(int));          *buffer = 1234;
    printf("Buffer address: %p\n",        &buffer);
    printf("Buffer contents: %p\n",      buffer);
    printf("Dereferenced value: %d\n",    *buffer);
    return 0;
}
```

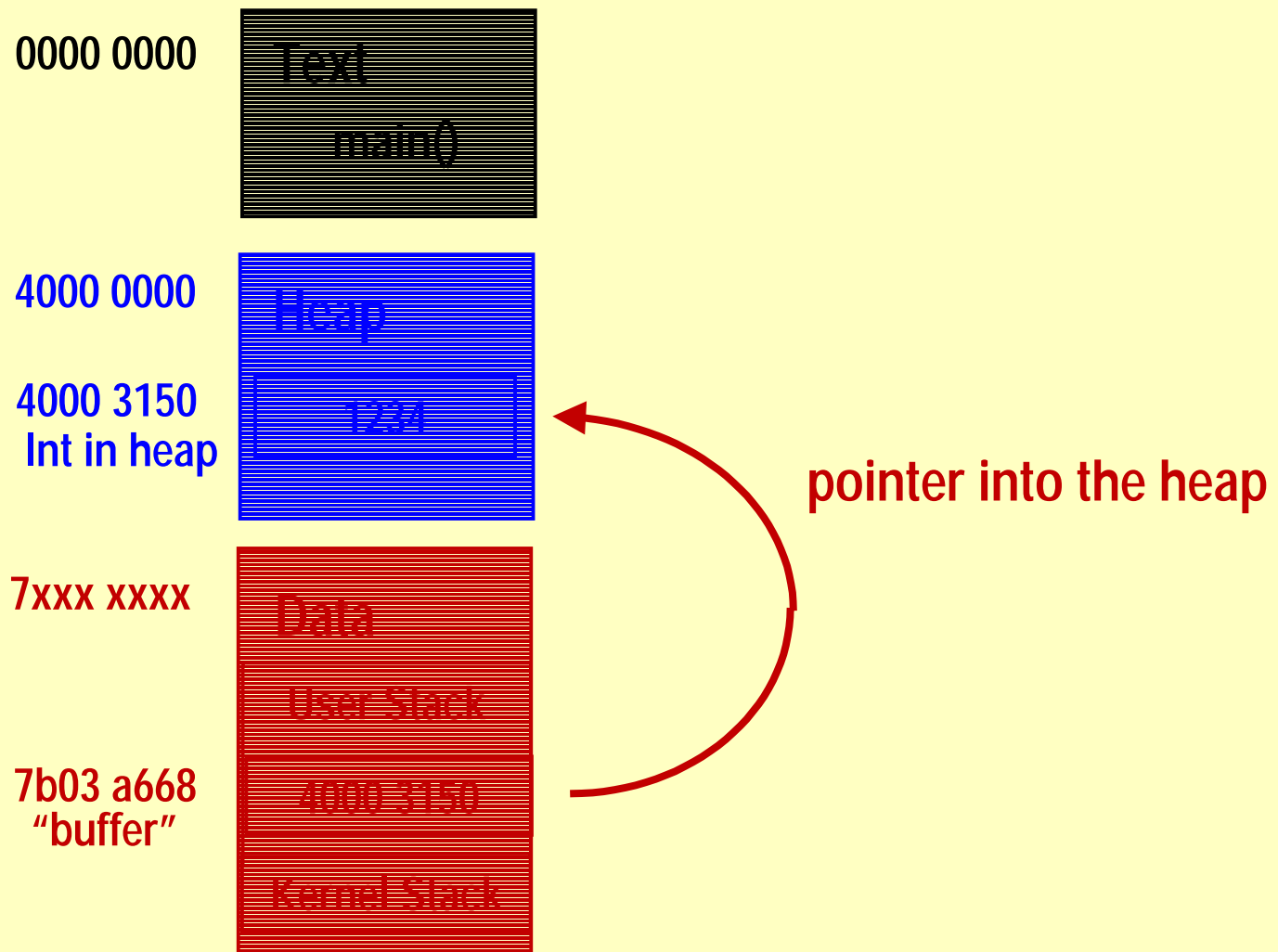
**HP-UX 10.20**

**ILP32**



```
$ ./malloc_return
Buffer address: 7b03a668  <-- address in data segment
Buffer contents: 40003150 <-- address in heap
Dereferenced value: 1234 <-- dereference ptr in heap
```

## Example Program - Memory Map



Example: code works on 32-bit  
and will fail on 64-bit

***DAL***

```
int main ()
{
    int *buffer;
    buffer = malloc(sizeof(int));
    *buffer = 1234;
    printf("Buffer address: %p\n",      &buffer);
    printf("Buffer contents: %p\n",    buffer);
    printf("Dereferenced value: %d\n", *buffer);
    return 0;
}
```

```
$ getconf KERNEL_BITS
64
$ cc +DD64 -o malloc_return
malloc_return.c
$ ./malloc_return
Memory fault(coredump)
```

**HP-UX 11.0**  
**LP64**



# Example: code works on 32-bit and will fail on 64-bit

***DAL***

```
int main ()
{
    int *buffer;
    buffer = malloc(sizeof(int));
    *buffer = 1234;
    printf("Buffer address: %p\n",      &buffer);
    printf("Buffer contents: %p\n",    buffer);
    printf("Dereferenced value: %d\n", *buffer);
    return 0;
}
```

- In C an undefined function return value is integer type 32 bits
- malloc returns a pointer type which is now 64-bits
- pointer.64 to integer.32 data truncation → invalid pointer value
- invalid pointer dereference → core dump



# Example: code works on 32-bit and will fail on 64-bit

***DAL***

```
#include <stdlib.h>
int main ()
{
    int *buffer;
    buffer = malloc(sizeof(int));
    *buffer = 1234;
    printf("Buffer address: %p\n", &buffer);
    printf("Buffer contents: %p\n", buffer);
    printf("Dereferenced value: %d\n", *buffer);
    return 0;
}
```

- Correct this by including the appropriate function declaration from the include file <stdlib.h>
- lint would have detected this missing function declaration

# Example: code works on 32-bit and will fail on 64-bit

**DAL**


```
#include <stdlib.h>
int main ()
{
    int *buffer;
    buffer = malloc(sizeof(int));
    *buffer = 1234;
    printf("Buffer address: %p\n",
          &buffer);
    printf("Buffer contents: %p\n",
          buffer);
    printf("Dereferenced value: %d\n",
          *buffer);

    return 0;
}
```

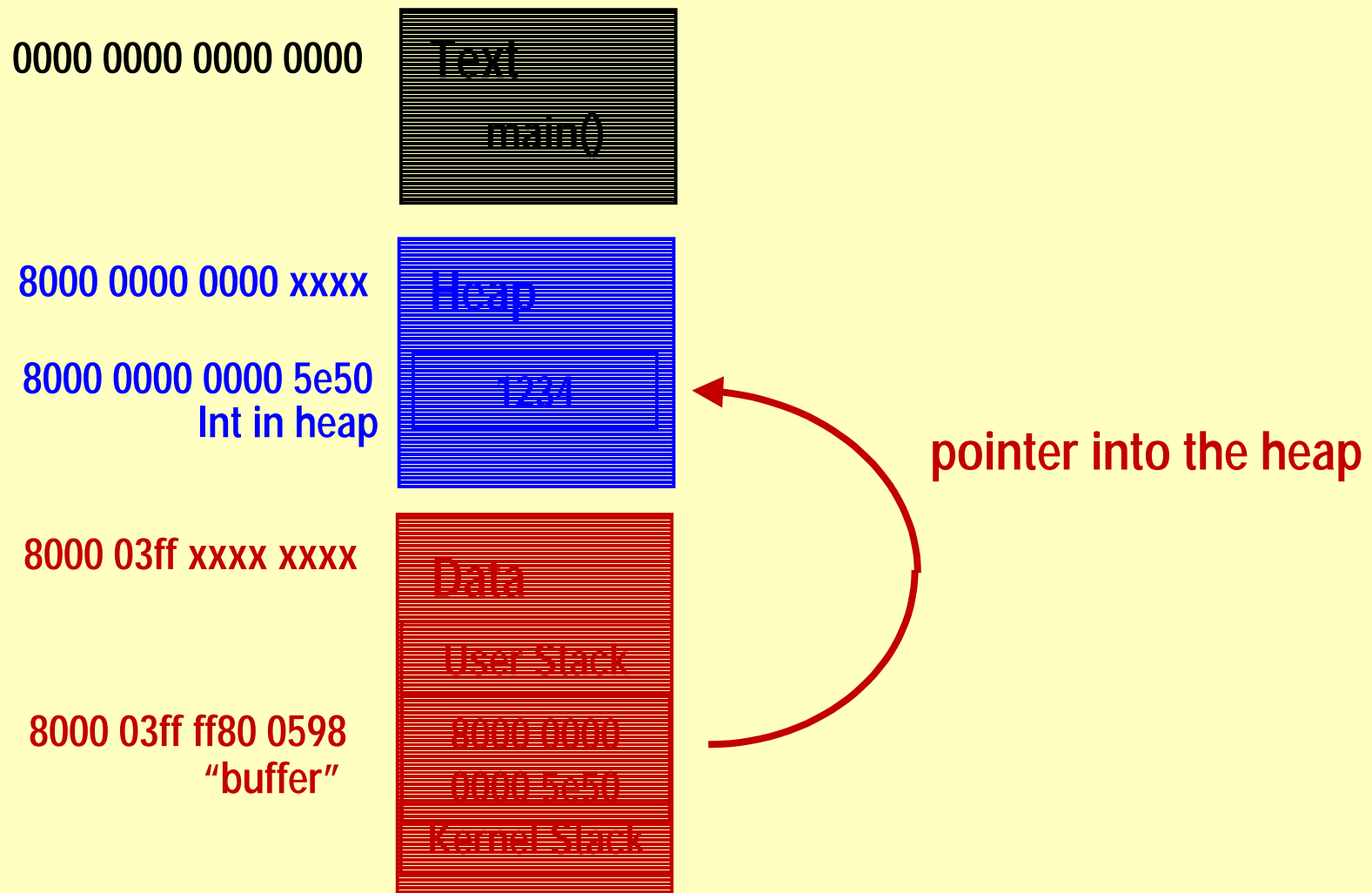
```
$ cc +DD64 -o malloc_return2 malloc_return2.c
$ ./malloc_return2
```

```
Buffer address: 800003ffff800598 <-- data segment
Buffer contents: 8000000000005e50 <-- ptr in heap
Dereferenced value: 1234 <-- dereference ptr.
```

**HP-UX 11.0  
LP64**



## Example Program - Memory Map



# Example: compiling and running on 32- and 64-bit systems

***DAL***

```
/*
 * test_basetypes: display maximum values of base data types
 */

#include <values.h>

main()
{
    printf("** Maximum values of base types on this system **\n");
    printf("Max int   = %ld\n", MAXINT );
    printf("Max long = %ld\n", MAXLONG);
}
```

# Example: compiling and running on 32- and 64-bit systems

***DAL***

```
$ uname -a  
HP-UX te887-01 B.11.00 A 9000/887 312087281 two-user license
```

```
$ getconf KERNEL_BITS  
32
```

```
$ cc -o test_basetypes test_basetypes.c
```

<-default compile: 32-bit

```
$ file test_basetypes  
test_basetypes: PA-RISC1.1 shared executable dynamically linked -not stripped
```

```
$ ./test_basetypes  
** Maximum values of base types on this system **  
Max int = 2147483647  
Max long = 2147483647
```

```
$ cc -o test_basetypes +DD64 test_basetypes.c
```

<- compile for 64-bit

```
$ file test_basetypes  
test_basetypes: ELF-64 executable object file - PA-RISC 2.0 (LP64)
```

```
$ ./test_basetypes  
sh: ./test_basetypes: Execute permission denied.
```

<- cannot execute 64-bit application  
on 32-bit system

# Example: compiling and running on 32- and 64-bit systems (continued) *DAL*

```
$ getconf KERNEL_BITS  
64
```

```
$ cc -o test_basetypes test_basetypes.c <-default compile: PA 2.0 32-bit  
/usr/ccs/bin/ld: (Warning) At least one PA 2.0 object file (test_basetypes.o)  
was detected. The linked output may not run on a PA 1.x system.
```

```
$ ./test_basetypes <-run the PA 2.0 32-bit application  
** Maximum values of base types on this system **  
Max int = 2147483647  
Max long = 2147483647
```

```
$ cc -o test_basetypes +DD64 test_basetypes.c <-compile for PA 2.0 64-bit
```

```
$ file test_basetypes  
test_basetypes: ELF-64 executable object file - PA-RISC 2.0 (LP64)
```

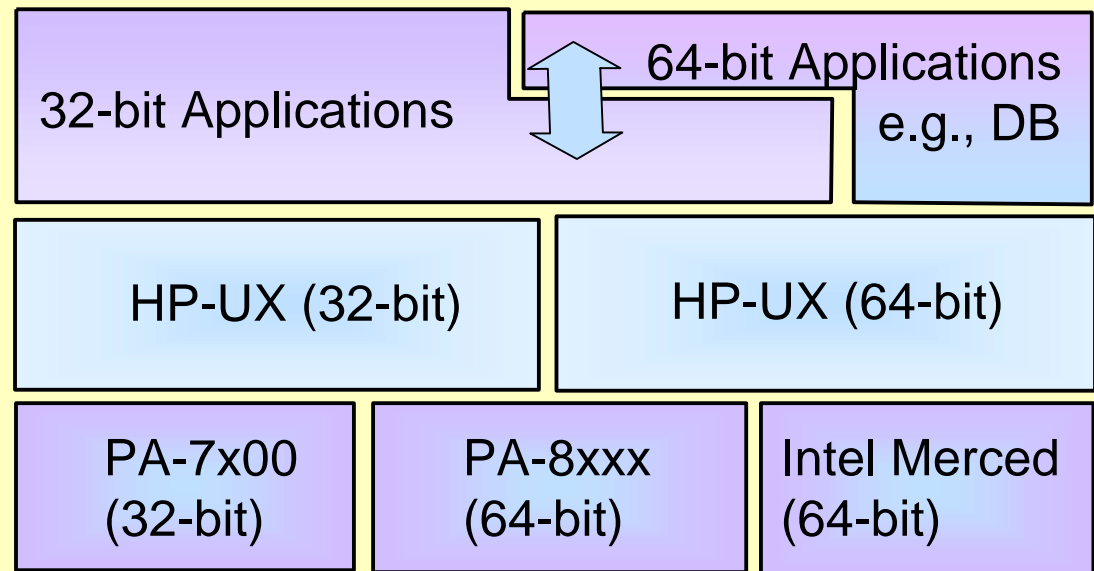
```
$ ./test_basetypes <-run the PA 2.0 64-bit application  
** Maximum values of base types on this system **  
Max int = 2147483647  
Max long = 9223372036854775807
```

```
$ cc -o test_basetypes +DA1.1 test_basetypes.c <-compile for PA 1.1 64-bit
```

# Smooth Upgrade to 64-bit Environment

***DAL***

- End-users may recompile and/or recompile their 32-bit applications at their leisure to access 64-bit functionality directly, or ...
- Run existing 32-bit applications unchanged and access 64-bit functionality of 64-bit databases or other ISVs



**HP-UX will allow co-existence and interoperability of 32- and 64-bit applications**

---

# HP-UX 11.0/10.30: Major New Functionality

***DAL***

---

- Merged binaries and merged install media
- Kernel threads
- Performance optimized page size
- Dynamically loadable kernel modules
- Fast core dump/recovery
- Fast password lookup
- Ignite UX
- New system recovery capability
- Year 2000 safety
- Other new technologies



---

# HP-UX 11.0 Feature Set

## New and Obsolete

***DAL***

---

- For a complete reference and pointers to other documentation, please browse the HP-UX Web Site:
  - <http://www.enterprisecomputing.hp.com/>
  - <http://www.unixsolutions.hp.com/>
- Other relevant web sites:
  - <http://docs.hp.com> (<http://www.docs.hp.com>)
  - <http://software.hp.com> (<http://www.software.hp.com>)

# Merged binaries and install media

***DAL***

- Merged install media
  - Install media now common for both S800 and S700
  - only 1 set of CD-ROMs
- Most S800/S700 dependencies removed from core OS
  - optional products (e.g., graphics) may still vary
- Merged install kernels
  - cold install kernel is the same for both S800 and S700

# Merged binaries and install media (continued)

**DAL**

- Merged binaries in core OS for 32- and 64-bit systems
  - every command in the core OS is built in 32-bit mode and can operate on either a 32- or 64-bit system
  - some kernel-specific commands are both 32- and 64-bit capable (such as ps and dm esg)
- Machine IDs and specifiers
  - future machine IDs are to be S700/S800 neutral
    - remove any dependency on S700/S800 machine id
  - New V-class machine types:

	uname -m	model
V2200	9000/800	9000/800/V2200
V2500	9000/800	9000/800/V2500

# New ELF object file format for 64-bit executables

***DAL***

- For 32-bit applications, the object file format remains SOM-32 (Spectrum Object Module).
- For 64-bit applications, the object file format is now ELF-64 (Executable and Linking Format)
  - ELF: a binary format developed by UNIX System Laboratories, used on Solaris, Linux, others.
  - New executable file startup mechanism
  - More efficient executable file format
- Why keep SOM for 32-bit and adopt ELF for 64-bit?
  - Maintain compatibility existing (32-bit) code
  - gain benefits for newly developed (64-bit) code

---

# HP-UX Threads in the 10.X Releases

***DAL***

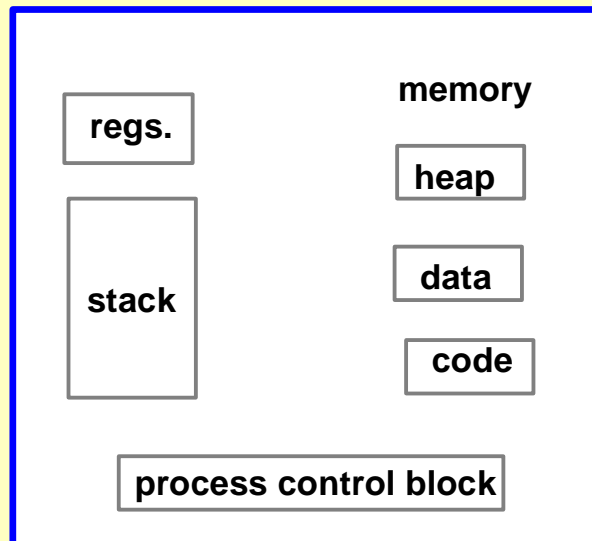
---

- HP-UX 9.X and 10.X offered User Space threads
- All Thread APIs conform to POSIX standards
  - (Pthread standard)
- Threads libraries, include files, and documentation delivered with the core DCE product
- Implementation based on OSF DCE Threads
- Actual thread libraries are contained in libcm a
  - CMA = Concert Multithreaded Architecture

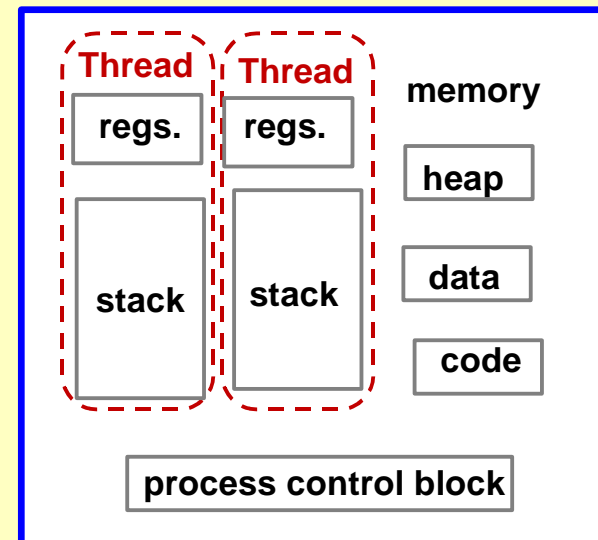
## Thread Definition

- A single sequential flow of control that can coexist with other threads in the same process.
- A thread is described by:
  - a unique identifier - Thread ID (TID)
  - shared address space with other threads in process
  - scheduling priority and policy

### Traditional Process



### Multi-threaded Process

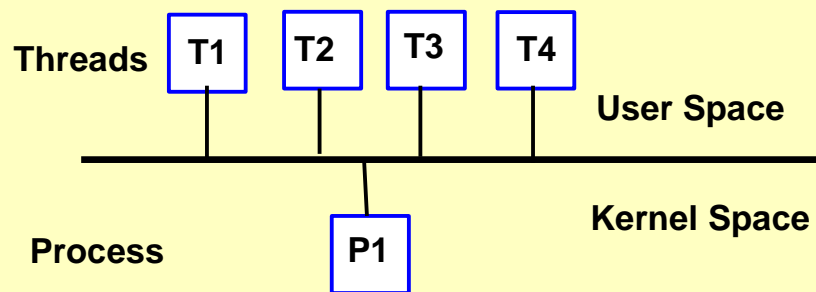


# User Threads and Kernel Threads

**DAL**

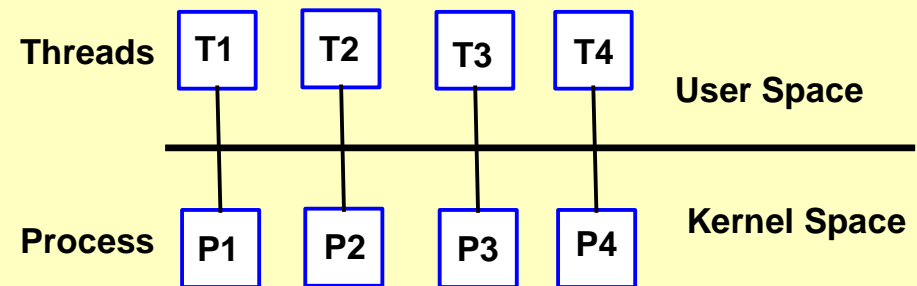
## User Threads

- handled in User Space and controlled using the threads APIs provided in the threads libraries
- M x 1 : Many to One Model. M Threads exist inside 1 Process that is visible to the Kernel



## Kernel Threads

- handled in User and Kernel space and are kernel schedulable entities visible to the operating system
- 1 x 1 : One to One Model. one Thread in one Kernel process. HP-UX implements with LWP (Light Weight Process).



# User Threads vs. Kernel Threads

***DAL***

## M x 1 User Threads

- all threads mapped into one process
- not visible to kernel
- fast creation/control
- no kernel overhead
- if one thread blocks, all threads are blocked

## 1 x 1 Kernel Threads

- each user thread mapped to kernel thread
- each thread scheduled independently, so if one blocks others can still run
- some overhead due to kernel calls



---

# Kernel Threads

## HP-UX Implementation

***DAL***

---

- Kernel thread infrastructure delivered in the 10.20 Release:
  - thread structure in kernel
  - kernel scheduling done on thread objects
- Full implementation of Kernel Threads delivered in 10.30 and 11.0 Releases

---

# Threads - Concurrency and Parallelism

***DAL***

---

- Concurrency - multiple threads may be in progress at the same time
- Parallelism - multiple threads are executing at the same time
  - On a M P system , Kernel Threads may execute in parallel on different CPU s
  - potentially a huge performance win

# POSIX Pthreads

- Pthreads can be implemented with M x 1, 1 x 1, or M x N models
- POSIX Standard: P1003.1c Portable API for Threads Extensions to POSIX 1003.1(a,b)
- HP-UX User Space (DCE) threads comply with POSIX Draft 4.
- HP-UX Kernel threads (HP-UX 11.0) comply with POSIX Draft 10.

# Programming with Threads

- Write programs to create and execute simple multi-threaded processes.
- Will demonstrate the basics of thread programming
- Will demonstrate Multi-Processor scaling and concurrency issues
- Test program will be a multi-threaded: "Hello, World!", then will become a form of compute-bound application

# Programming with Threads

## program pseudo-code

*DAL*

```
/*
 * Create and execute 2 threads (pseudo-code)
 */

main()
{
  pthread_create(thread1);          /* create first thread */
  pthread_create(thread2);         /* create second thread */

  pthread_join(thread1);           /* wait for first thread */
  pthread_join(thread2);           /* wait for second thread */
}

thread1()                           /* Thread 1 will run here */
{
  printf("Hello, World!\n");
}

thread2()                           /* Thread 2 will run here */
{
  printf("Goodbye, World!\n");
}
```

# Programming with Threads - program code

*DAL*

```
/*
 * Create and execute 2 threads (main program)
 */

#include <pthread.h>
void thread1(); /* forward declarations */
void thread2(); /* for thread start addr */

main()
{
    pthread_t tid1,tid2; /* Thread ID values */

    /* Create threads */
    pthread_create(      &tid1, /* Thread ID (returned value) */
        (pthread_attr_t *) NULL, /* thread attributes */
        (void *(*)(void *)) thread1, /* thread executes this func */
        (void *) NULL ); /* argument passed to thread */
    pthread_create(      &tid2,
        (pthread_attr_t *) NULL,
        (void *(*)(void *)) thread2,
        (void *) NULL );

    /* wait for the threads to finish */
    pthread_join(        tid1, /* wait on this thread */
        (void **) NULL ); /* thread's exit value */
    pthread_join(        tid2,
        (void **) NULL );
}
```

# Programming with Threads - program code (continued)

***DAL***

```
/*  
 * Functions where the two threads will start executing  
 */
```

```
void thread1()  
{  
    printf("Hello, World!\n");  
    pthread_exit( (void *) NULL);  
}
```

```
/* return a value and exit */
```

```
void thread2()  
{  
    printf("Goodbye, World!\n");  
    pthread_exit( (void *) NULL);  
}
```

# Programming with Threads - compiling and executing

***DAL***

- compile the program

```
% cc +DA1.1 -lpthread -o 2threads 2threads.c
```

- compile +DA1.1 so the binary will be portable and will run on any supported 11.0 system

- Pthreads library located at /usr/lib/libpthread.[a,sl]

(for DCE threads link against libcma)

- default compilation builds a shared-linked executable

- run the program

```
% ./2threads
```

```
Hello, World!
```

```
Goodbye, World!
```



# Programming with Threads - order of execution

***DAL***

- run the program

```
% ./2threads
```

```
Hello, World!
```

```
Goodbye, World!
```

*... OR MAYBE ...*

```
% ./2threads
```

```
Goodbye, World!
```

```
Hello, World!
```

- Is the order of execution guaranteed?

It is not guaranteed unless one uses the thread scheduling and synchronization APIs

---

# Programming with Threads - compute-bound example

***DAL***

---

- Each thread will now execute a basic example of a compute-bound application
- Threads will execute concurrently as the kernel schedules each thread
- Threads may execute in parallel on a multi-processor system
- Default thread scheduling policy on HP-UX is to schedule threads across different processors.

# Programming with Threads - compute-bound program code

**DAL**

```
#include <pthread.h>
void busywork(); /* where the threads execute */

main()
{
    pthread_t tid1,tid2;

    /* Create threads */
    pthread_create( &tid1, /* Thread ID (returned value) */
        (pthread_attr_t *) NULL, /* thread attributes */
        (void *(*)(void *)) busywork, /* thread executes busywork */
        (void *) 10000 ); /* argument passed to thread */

    pthread_create( &tid2,
        (pthread_attr_t *) NULL,
        (void *(*)(void *)) busywork,
        (void *) 10000 );

    /* wait for the threads to finish */
    pthread_join( tid1, (void **) NULL); }
    pthread_join( tid2, (void **) NULL);
}
```

```
void busywork(int maxcount)
{
    int count1, count2;
    for (count1=0; count1<maxcount; count1++)
        for (count2=0; count2<maxcount; count2++);
    pthread_exit( (void *) NULL);
}
```

**busywork: loop for  
(count\*count) times**

# Programming with Threads - compute-bound execution

***DAL***

- Program compiled the same as before

```
% cc +DA1.1 -lpthread -o 2busythreads 2busythreads.c
```

- execute the program on 1-way C-110:

```
% timex 2busythreads
```

```
real      17.10
```

```
user      16.79
```

```
sys       0.03
```

- execute the program on 2-way D-270:

```
% timex 2busythreads
```

```
real      5.06
```

```
user     10.05
```

```
sys       0.01
```

Real (wall clock) time

Compute time 2X real time!

---

# Program m ing w ith Threads - com pute-bound exam ple

***DAL***

---

- Real tim e (w all clock tim e) is halved on the 2-w ay system
  - Each thread executes in parallel on a different processor
- This program M P-scales perfectly to a 2-w ay system , but w ould not scale any better on a system w ith m ore than 2 processors installed.
- This sim ple program exam ple is not concerned w ith synchronization , global resources and contention , asynchronous events -m any of the interesting elem ents of real-w orld program m ing .

# Threads References

- Thread Time, Scott Norton and Mark DiPasquale, HP Press and Prentice-Hall, ISBN 013-190067-6
- Multi-Threaded Programming under HP-UX, A. Jyothy R. Ready, HP World August 1996
- HP-UX Whitepaper in the documentation directory  
`/usr/share/doc/proc_mgt.[ps,txt]`
- FAQ in `comp.programming.threads`

# Large VM Page Size

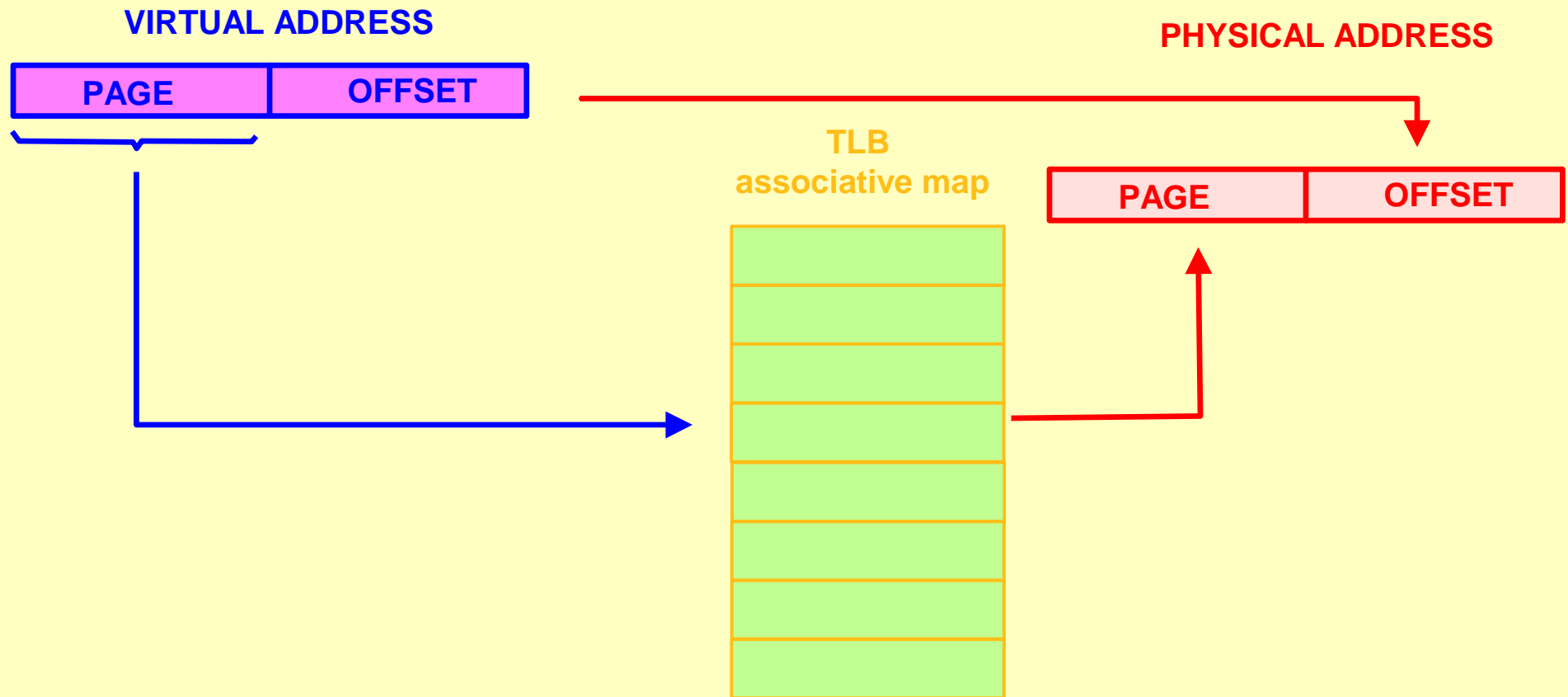
- Virtual Memory Page Size was fixed at 4 K Bytes for all HP-UX releases prior to 10.20
- HP-UX 10.20 introduced Large Pages - VM page size for an executable process could be explicitly set using the `chattr` command to:

4K	1M	256M
16K	4M	L (use largest available size)
64K	16M	D (use kernel default size)
256K	64M	

- Page size could be set for Text, Data, Stack, Memory Mapped Files, and other memory objects.
- Only for PA-8000 processors

# Virtual Memory Page Lookups and Translation Lookaside Buffer (TLB)

*DAL*



TLB is a hardware lookup (extremely fast) of virtual page addresses. If page entry is not in TLB, a software search (slower) of all virtual page addresses must be done by the kernel to get the physical page address.



# Performance Optimized Page Size

***DAL***

- HP-UX 11.0 introduces variable page size (aka. Large Pages or Performance Optimized Page Size).
- The SysAdmin can request a page size using the `chattr()` command and or the kernel can specify a page size.
- The kernel tries to honor the request but may use a smaller page size if there is competition from memory.
- Kernel tunable parameters to control operation:
  - `vps_pagesize` default page size used by kernel
  - `vps_ceiling` maximum page size used by kernel
  - `vps_chattr_ceiling` maximum size a user can set

*(probably don't want or need to change these tunables)*

---

# Advantages to Performance Optimized Page Size

***DAL***

---

- Larger virtual address ranges can be mapped using fewer TLB entries, so there will be fewer TLB misses
- PO PS will offer performance advantages for applications that:
  - are experiencing significant TLB misses
  - have large Reference Sets (e.g. large Data Segments or Text Segments)

# Performance Optimized Page Size Tools and Documentation

***DAL***

- `chattr` command (see `chattr(1)` man page)
  - `chattr+pi <size>` text page size
  - `chattr+pd <size>` data page size
  - size = 4K , 16K , 64K , 256K , 1M , 4M , 16M , 64M , 256M , L (largest), D (kernel default)
- kernel tunable parameters:
  - `vps_pagesize` default page size kernel will use
  - `vps_ceiling` max page size kernel will select
  - `vps_chattr_ceiling` max page size a user can select
- `/usr/contrib/bin/vps_stats`
  - report page statistics
- Whitepaper: `/usr/share/doc/mem_mgmt.[ps,txt]`

# Performance Optimized Page Size Example Programs

***DAL***

```
/*  
 * Allocate/touch large data space - good locality  
 */
```

```
#define SIZE 256
```

```
int big_array [SIZE] [SIZE] [SIZE];
```

← 65 MB Data Segment  
16 Million 4-byte ints

```
main( )
```

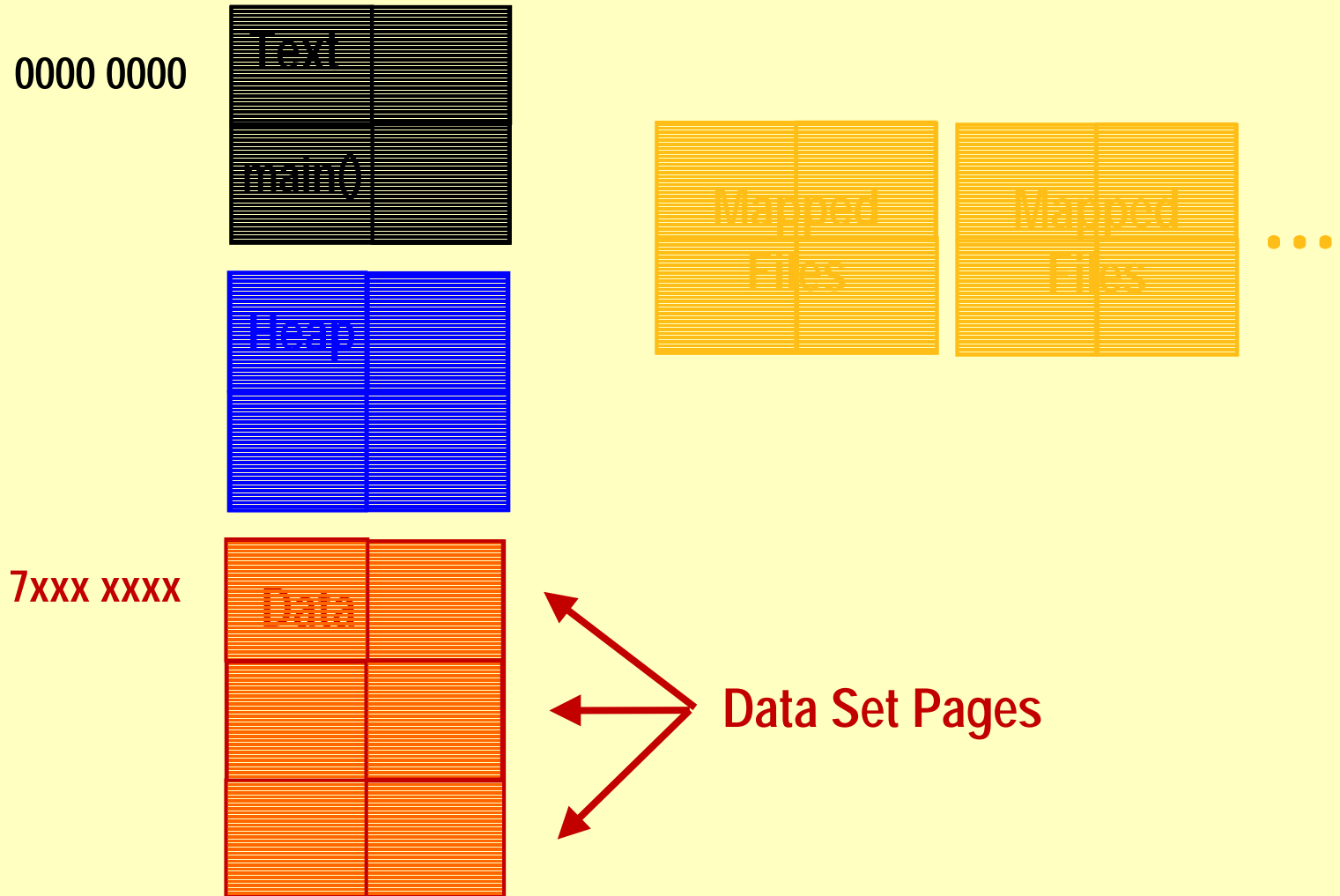
```
{  
  int i,j,k;      /* array indices */
```

```
  for ( i=0; i<SIZE; i++ )  
    for ( j=0; j<SIZE; j++ )  
      for ( k=0; k<SIZE; k++ )  
        big_array[i][j][k] = k;
```

← Touch each element  
16 Million iterations

# Memory Map - Virtual Memory Pages

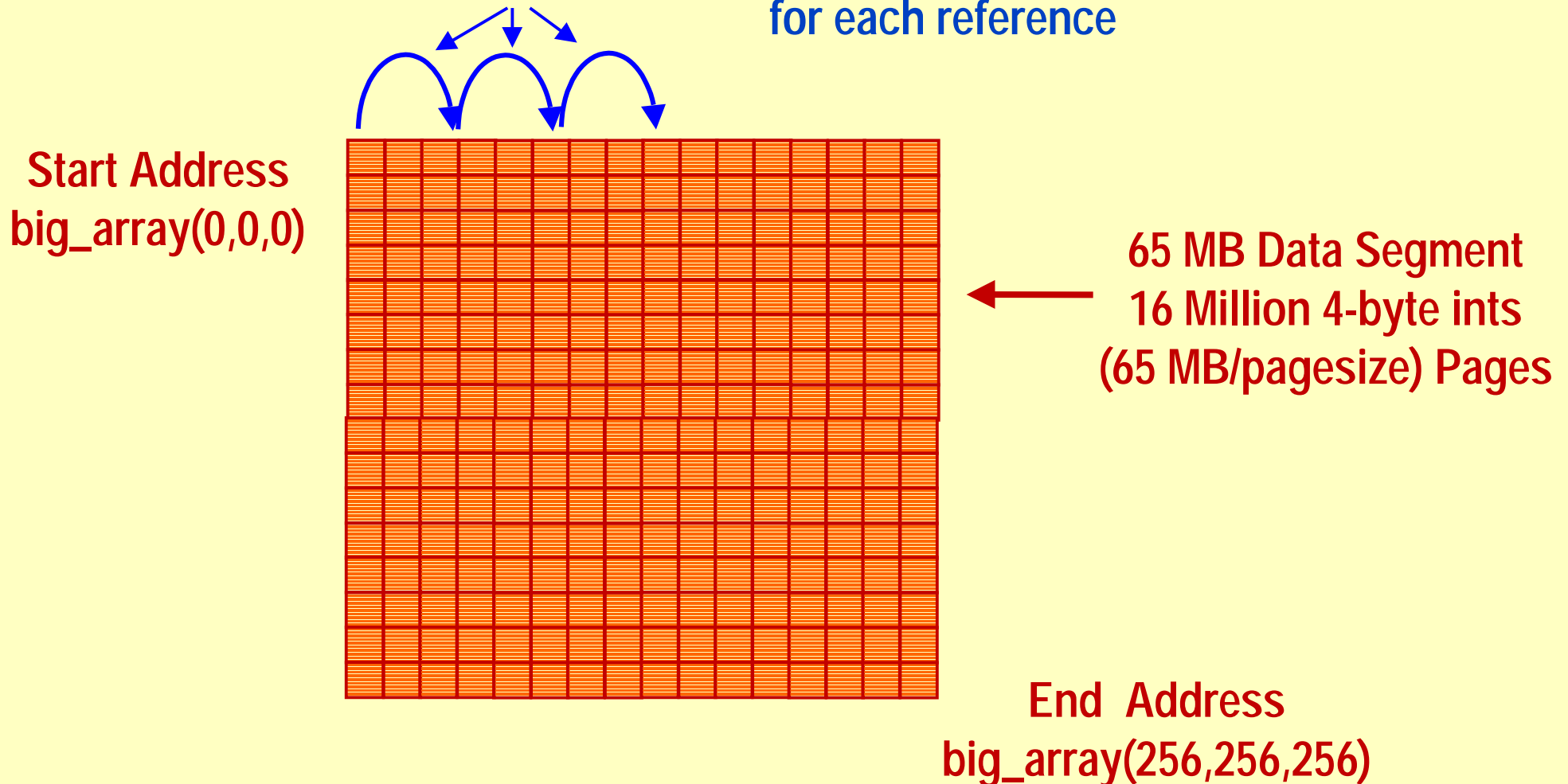
***DAL***



# Example Program - Good Locality

***DAL***

GOOD Data Locality: 4 byte address jumps  
for each reference



# Example Program - PoorLocality

**DAL**

```
/*  
 * Allocate/touch large data space - poor locality  
 */
```

```
#define SIZE 256
```

```
int big_array [SIZE] [SIZE] [SIZE];
```

← 65 MB Data Segment  
16 Million 4-byte ints

```
main()  
{
```

```
    int i,j,k;    /* array indices */
```

```
    for ( i=0; i<SIZE; i++ )
```

```
        for ( j=0; j<SIZE; j++ )
```

```
            for ( k=0; k<SIZE; k++ )
```

```
                big_array[k][j][i] = k;
```

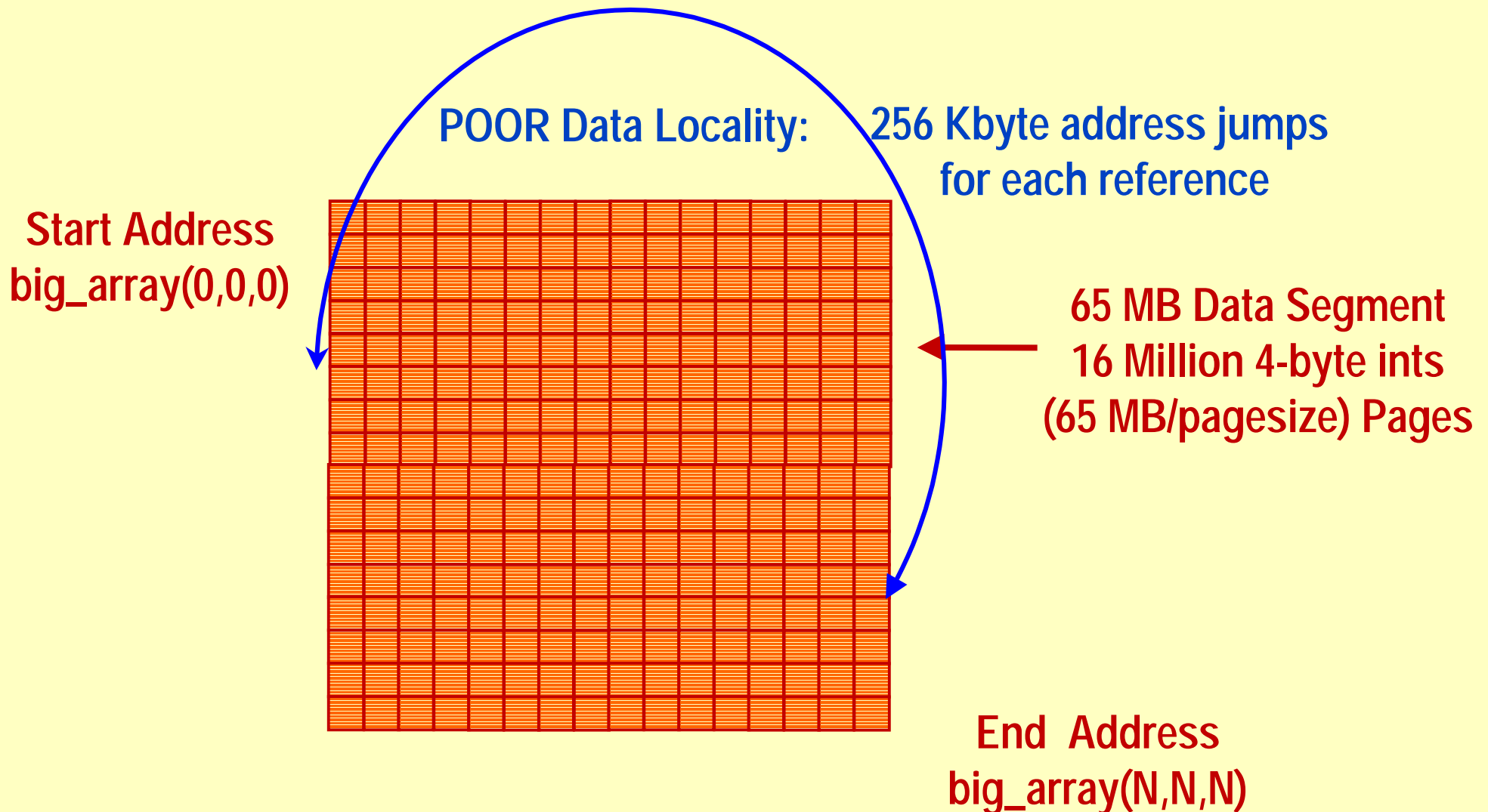
```
}
```

← Poor Locality -  
256K address jumps  
for each reference

*All we've done is reverse these indices!*

# Example Program - Poor Locality

***DAL***





---

# Test Program

## Default Compile and Run

***DAL***

---

- the default compile will set the page size to the kernel default page size:
  - chatr shows "D" for page size
- the default kernel tune is for 16K page size
- the Program Memory Map will consist of:
  - Text Segment: very small
  - Data Segment:  
4096 entries of 16 K byte pages = 65 M Bytes

# Test Program Memory Map:

**DAL**

```
vps_stats -p <process_id>
```

## REGION INFO

TYPE	HINT	PHYSPAGES	Phys Page SIZE[Count]
----	----	-----	-----
NULL	4K	1	
TEXT	4K	3	4K[3]
DATA	16K	16387	4K[3] 16K[4096]
MMF	4K	0	
MMF	4K	3	4K[3]
MMF	4K	1	4K[1]
MMF	4K	3	4K[3]
MMF	4K	4	4K[4]
MMF	4K	1	4K[1]
MMF	4K	2	4K[2]
MMF	4K	3	4K[3]
STACK	16K	3	4K[3]
MMF	4K	15	4K[15]
MMF	4K	206	4K[206]
MMF	4K	1	4K[1]
UAREA	4K	7	4K[7]

Data Segment:  
16 KB pages  
4096 entries  
= 65 Mbytes space

# Execution Times

## Default (16 KB) pages

*DAL*

- Execution time: measure with "time big\_array"
- TLB misses: measure with "cyclemeter" (contributed tool) (PerfUX or GlanceUX also work)

- Good Locality program :

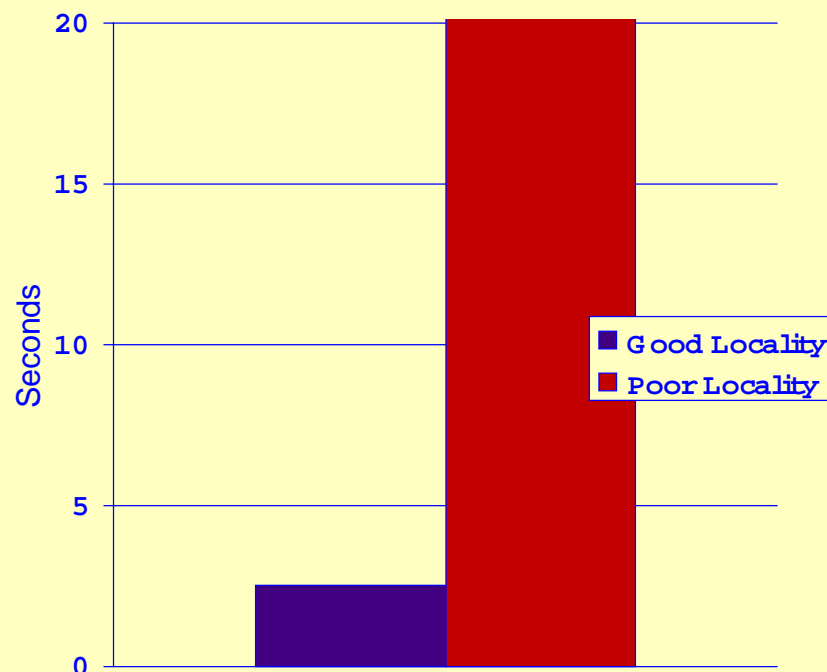
Runtime 2.52 sec.

Data TLB misses 51,535

- Poor Locality program :

Runtime 20.49 sec.

Data TLB misses 17,062,436



# Memory Map:

## char to 256K Data Pages

**DAL**

### REGION INFO

TYPE	HINT	PHYSPAGES	Phys Page SIZE[Count]
----	----	-----	-----
NULL	4K	1	
TEXT	4K	3	4K[3]
DATA	256K	16387	4K[3] 16K[4] 64K[3] 256K[255]
MMF	256K	0	
MMF	256K	3	4K[3]
MMF	4K	1	4K[1]
MMF	256K	3	4K[3]
MMF	4K	4	4K[4]
MMF	256K	1	4K[1]
MMF	256K	2	4K[2]
MMF	4K	3	4K[3]
STACK	256K	3	4K[3]
MMF	4K	15	4K[15]
MMF	4K	206	4K[206]
MMF	4K	1	4K[1]
UAREA	4K	7	4K[7]

Data Segment:  
 256 KB pages  
 255 entries +  
 1 fragmented entry  
 = 65 Mbytes space

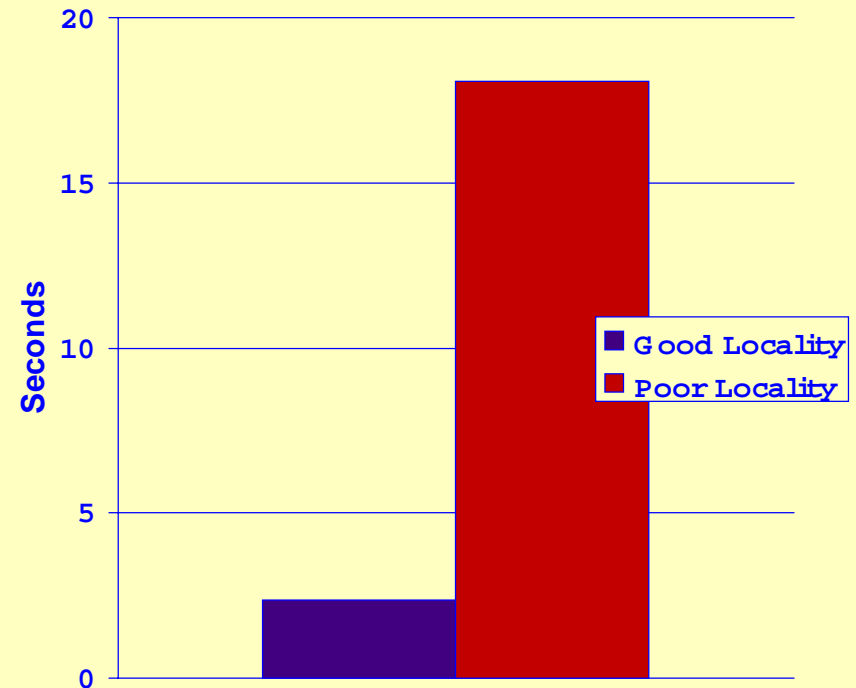


# Execution Times

## chatr to 256 K Byte pages

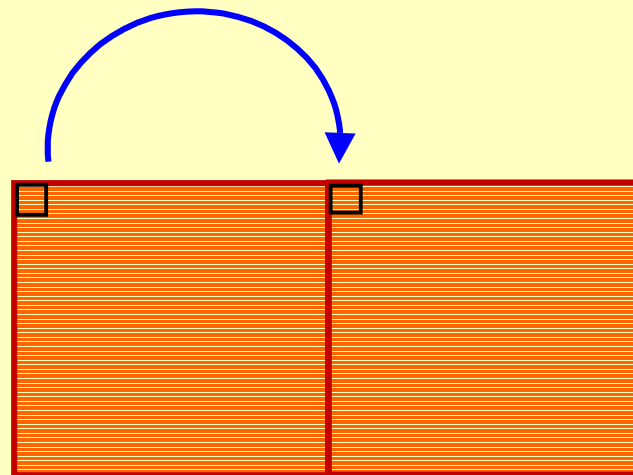
***DAL***

- Change Data Page size:  
"chatr+pd 256K big\_array"
- Good Locality program :  
Runtime 2.35 sec.  
Data TLB misses` 34,562
- Poor Locality program :  
Runtime 18.06 sec.  
Data TLB misses 16,906,509



## Change to 256 K byte Pages – why little performance change?

- Data Page Size is 256 K bytes
  - "PoorLocality" C-program references 4-byte integers from address (A B C) to (A + 65536 B C), ...
  - CurrentAddress + (4 bytes \* 65536) =  
CurrentAddress + 256K
- this is *justover* the next page boundary



# Memory Map:

# DAL

## chart to 4M Data Pages

### REGION INFO

TYPE	HINT	PHYSPAGES	Phys Page SIZE[Count]
----	----	-----	-----
NULL	4K	1	
TEXT	4K	3	4K[3]
DATA	4M	16387	4K[3] 16K[4] 64K[3] 256K[3] 1M[3] 4M[15]
MMF	4M	0	
MMF	4M	3	4K[3]
MMF	4K	1	4K[1]
MMF	4M	3	4K[3]
MMF	4K	4	4K[4]
MMF	4M	1	4K[1]
MMF	4M	2	4K[2]
MMF	4K	3	4K[3]
STACK	4M	3	4K[3]
MMF	4K	15	4K[15]
MMF	4K	206	4K[206]
MMF	4K	1	4K[1]
UAREA	4K	7	4K[7]

Data Segment:  
4 MB pages  
15 entries +  
1 fragmented entry  
= 65 Mbytes space



# Execution Times

## chatr to 4 M Byte pages

***DAL***

- Change Data Page size:  
"chatr+pd 4M big\_array"

- Good Locality program :

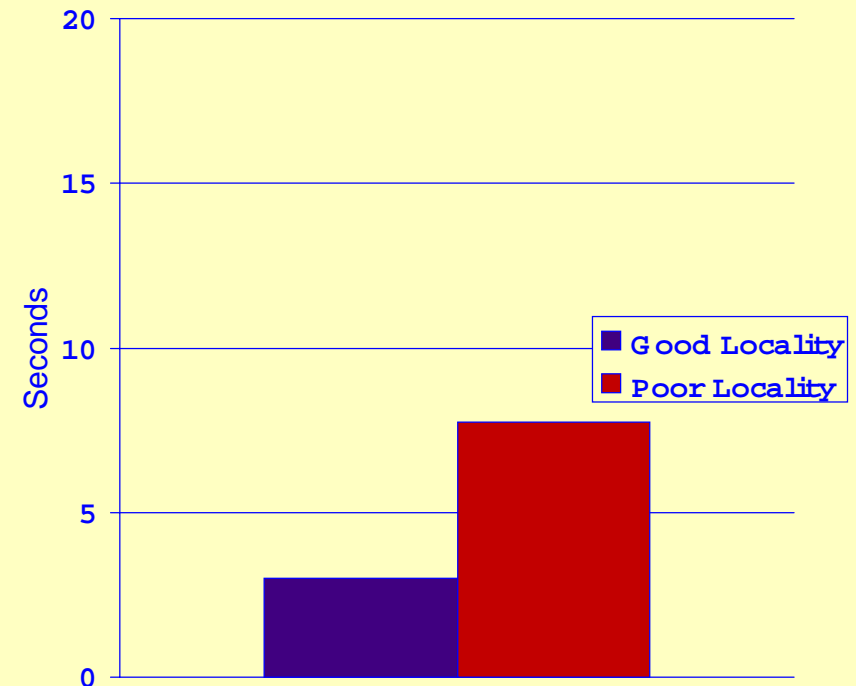
Runtime 2.99 sec.

Data TLB misses` 33672

- Poor Locality program :

Runtime 7.75 sec.

Data TLB misses 30,227





---

# Variable Pages - Other Considerations

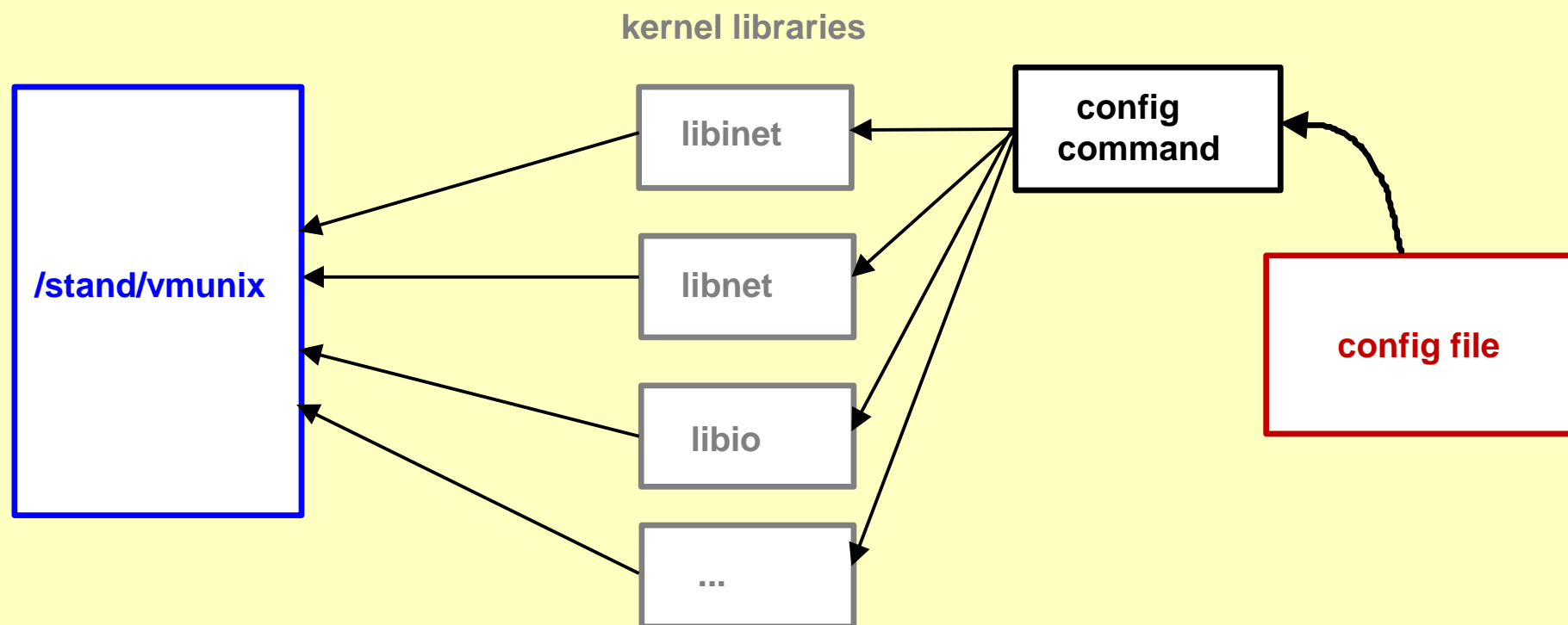
***DAL***

---

- Other factors affecting performance:
  - other ongoing system activity
  - memory page contention/thrashing/locking issues
  - cache sizes and cache fits
- Best performance is a combination of:
  - ✓ fast and efficient hardware (PA-8000)
  - ✓ correct system and kernel tunes
  - ✓ well-written software

## Static Kernel Configuration

- Prior to 11.0 Release, HP-UX kernels have been statically configured and built
- `config command` to build kernel

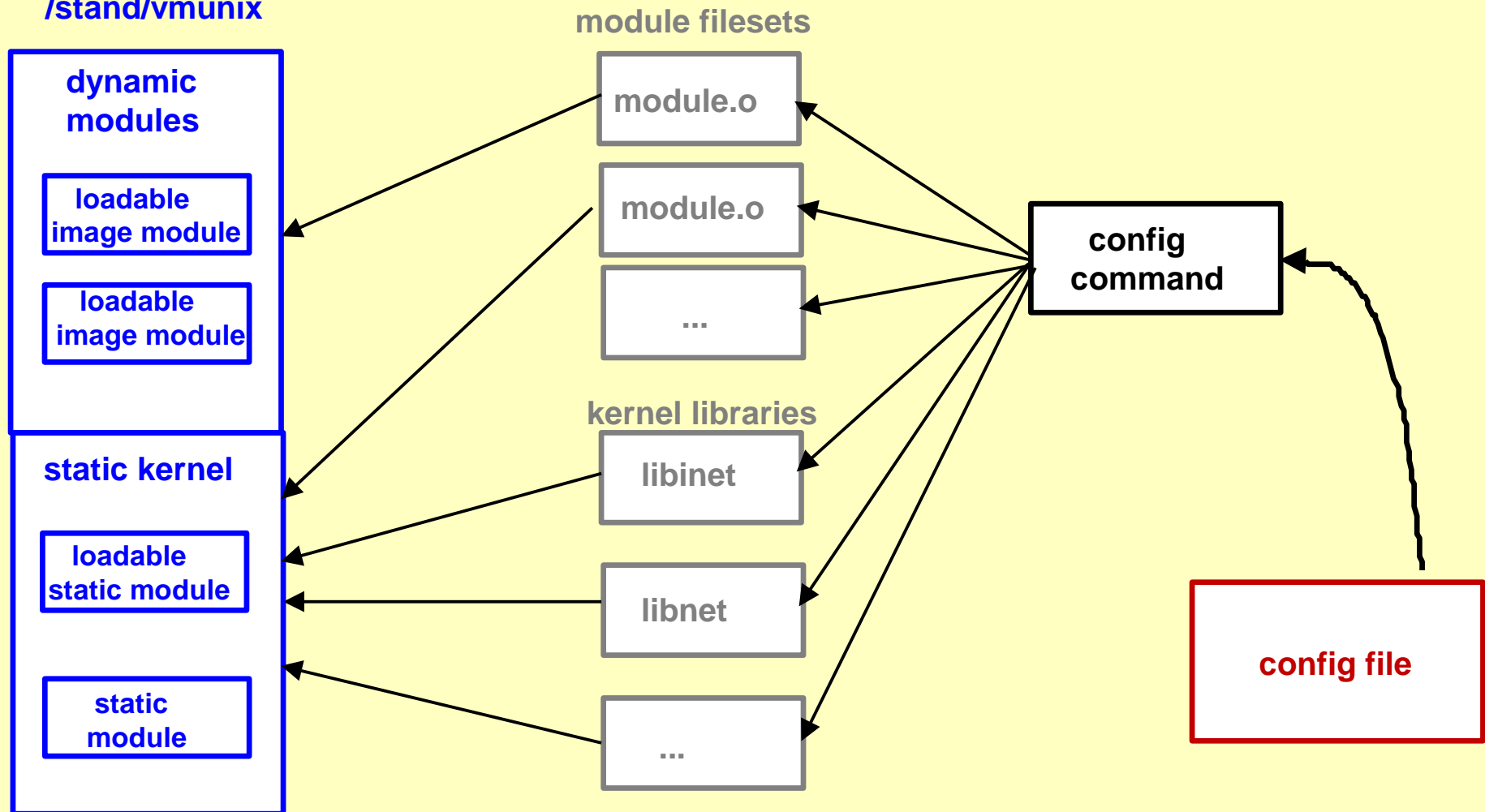


# Dynamic Kernel Configuration

**DAL**

Dynamically Loadable Kernel Modules allow one to load, unload, and configure kernel modules without rebooting

*/stand/vmunix*



# DLKM in 11.0

- DLKM Phase 1 delivered in 11.0:
  - DLKM infrastructure in kernel and commands
  - Major subsystems made DLKM-aware
  - no modules configured as dynamically loadable (yet)

# Configuring Kernels

- Static kernel configuration remains exactly the same
  - `config com m and and kernel config files`
- To build a new kernel module:
  - `config -M <module>`
- To update the kernel with the new module:
  - `config -u /stand/system -or-`
  - `km update (new com m and)`

---

# New Commands for configuring Kernels

***DAL***

---

<b>kmsystem</b>	<b>set control flags in system files</b>
<b>kmtune</b>	<b>interface to set tunable parameters</b>
<b>kmupdate</b>	<b>update system with new kernel or loadable modules</b>
<b>kminstall</b>	<b>install/remove/update a module in a system</b>
<b>kmmodreg</b>	<b>register/unregister a module with the system</b>
<b>kmadmin</b>	<b>general administrative interface for DLKM</b>

---

# D L K M – Future Capabilities

***DAL***

---

- On-Line Module Replacement
- Module support for more subsystems:
  - File systems
  - System calls
  - Other drivers

# Fast Core Dump/Recovery

- Support of Large Memory (to 16 G Bytes) means system image dump/recovery could take a long time
  - from several minutes to several hours
- Long dump/recovery times would not allow the High Uptime goals for HP-UX
  - one cannot afford to take several hours to dump/recover a system image
- Size of the dump image could be very large, and there may not be enough space to store the image



# Fast Core Dump/Recovery

New /renamed commands:

crashconf

Add/remove/redefine dump devices at runtime. Devices can also be statically declared with a new field in /etc/fstab. This can override the dump device that has been configured in the kernel.

savecrash  
(savecore)

Save system image during the next reboot. Can select full/partial, or compress/nocompress

crashutil

Formats the crash dump image for further analysis

# Fast Core Dump/Recovery

- Background mode operation of savecrash service at system startup
- There is a new mode flag to allow Background Operation available to startup scripts
  - Allows system startup to proceed while a lengthy process continues in background
  - Don't use this mode if other services depend on the immediate result of the startup of your service, or else the results may be confusing
- Background Mode startup will display "OK" in the startup screen if the service begins successfully

# Fast password/uid/gid Lookup

- Extension to uid/gid/password lookup for local files.
  - applies partially to NFS systems
  - does not apply to Secure systems
- Large numbers of uids (to 4 G .) means potentially large number of entries in /etc/passwd file.
- Existing APIs (getpwent(3c), getgrent(3c)) do a linear search through /etc/passwd - potentially long times required for simple login or lookup

# Fast Passwords

- Solution to lengthy searches of password files
  - a helper process to speed the lookup
- New system process: **pwgrd** - a user space daemon
- All supported APIs such as `getpwent(3C)` and `getgrent(3C)` communicate to **pwgrd**
- **pwgrd** hashes and caches password entries to avoid potentially lengthy searches of */etc/passwd*

# Fast Passwords

- **pwgrd** fast lookup
  - hashing of all entries in */etc/passwd* - done at startup of the daemon or if the daemon is signaled
  - caching of recent lookups into */etc/passwd*
- **pwgrd** caches recent retrievals from NIS entries
- Documentation:
  - **pwgrd** (1M) man page on 11.0 system

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# Ignite-U X

***DAL***

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- A new Cold Installation method
  - not an Update in Place method
- Available on 10 X (starting with DART 34) and 11.0

# Ignite-U X Capabilities

***DAL***

	<b>Cold Install</b>	<b>Ignite-U X</b>
Install methods	From local media Pull from network server	From local media Pull from network server Push from network server Reinstall existing machines
Server capability	Single Release	Multiple Releases
User Interface	TUI on target	TUI on target GUI on server
Software Formats	Software Distributor	Software Distributor, cpio, tar can use multiple formats

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# Ignite-U X

***DAL***

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- Install methods:
  - Boot from local media
  - Pull from network server
  - Push from network server
- uses `rem sh` to target system console
  - can easily add new clients



## Creating Custom Install Packages

- Customize an install configuration and save the configuration for replication
  - "Save As" option to Ignite-U X
- Ignite-U X tools allow one to create a custom archive (Golden Image) and choose the install format

# System Recovery

- Use to recover Root Filesystem with all your installation and customization
- Root recovery might ordinarily require these steps:
  - cold install
  - configure
  - reinstall patches
  - reinstall applications
  - reinstall user data and files

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# System Recovery - other uses

***DAL***

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- Modify RootFilesystem size
- Modify primary swap size
- Convert RootFilesystem from HFS to VxFS
- Clone a system

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# System Recovery

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- Supported on 11.0 and 10.X
- Installed with Ignite-UX
- Install from DART releases or HP Web site:
  - [http://www .softw are.hp.com /](http://www.softw are.hp.com/)
  - Link to: "Network and System Administration Tools"

# System Recovery

***DAL***

make_recovery	make a system recovery tape ("make_recovery -A" - entire core VG)
save_config	create a configuration file that details the current system hardware and software configuration
check_recovery	compare current system configuration to last configuration file
print_manifest	print hardware configuration (CPUs, LVM and disks, I/O), OS configuration, installed products)

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# System Recovery

***DAL***

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- Currently supports:
  - HFS/XFS
  - SAM interface
  - LVM
  - DDS tape recovery in age
- Not yet supported:
  - Striped or mirrored disks
  - Shared LVM

# System Recovery - print\_manifest utility

# DAL

## System Information

Your Hewlett-Packard 9000 computer has software installed and configured as follows.

-----  
NOTE: You should retain this information for future reference.  
-----

## System Hardware

Model: 9000/777/C110  
Main Memory: 128 MB  
Processors: 1  
OS mode: 32 bit  
HW capability: 32 bit  
LAN hardware ID: 0x0060B001BA5B  
Software ID: 2011905808  
Keyboard Language: PS2\_DIN\_US\_English

## HARDWARE CONFIGURATION

Storage devices	HW Path	Interface
SEAGATE ST15150W 4095 Mb	8/12.6.0	GSC built-in Fast/Wide SCSI Interface
SEAGATE ST32430N 2048 Mb	8/16/5.4.0	Built-in SCSI
HP C1533A	8/16/5.3.0	Built-in SCSI

## I/O Interfaces

Class	H/W Path	Driver	Description
ext_bus	8/12	c720	GSC built-in Fast/Wide SCSI Interface
ext_bus	8/16/0	CentIf	Built-in Parallel Interface
audio	8/16/1	audio	Built-in Audio
tty	8/16/4	asio0	Built-in RS-232C
ext_bus	8/16/5	c720	Built-in SCSI
lan	8/16/6	lan2	Built-in LAN
ps2	8/16/7	ps2	Built-in Keyboard/Mouse
pc	8/16/10	fdc	Built-in Floppy Drive
hil	8/20/1	hil	Built-in HIL
tty	8/20/2	asio0	Built-in RS-232C
graphics	10/16	graph3	Graphics

## Installed Software

Your system was installed with HP-UX version B.11.00.

Your system has the following software products installed and configured on the system disk drive(s).

Product	Revision	Description
B3782EA	B.10.20	HP-UX Media Kit (Reference Only. See Description)
B3884EA_AGN	B.10.20	HP-UX 32-User License
B3899BA	B.11.01.01	HP C/ANSI C Developer's Bundle for HP-UX 11.00 (S700)
B3911DB	B.11.01.01	HP aC++ Compiler (S700)
B3919EA_AGS	B.11.00	HP-UX Unlimited-User License
B4580AA	B.11.00.01	HP-UX 11.00 Software Transition Kit
B5455CA	C.01.16.00	HP-UX Development Kit for Java*
B5724AA_APZ	A.1.45	HP-UX Installation Utilities (Ignite-UX - S700 - 10.20)
DCEProg	B.10.20	DCE Programming and Archive Libraries
HPUXEngCR700	B.10.20	English HP-UX CDE Runtime Environment
J2559C	D.06.15	Hewlett-Packard JetAdmin for Unix Utility
UXCoreMedia-J	B.11.00	HP-UX Japanese Media Kit (Reference Only. See Description)
XSWGRI100	B.11.00.39	HP-UX Extension Pack, June 1998

## INSTALLED SOFTWARE

# System Recovery - print\_m anifest (continued)

**DAL**

## LVM File System Configuration

This system is configured with Logical Volume Manager (LVM) file systems.  
Refer to the File System layout section for information on the LVM layout.

## Disk layout

LVM disk	Device file	HW Addr	size	vol. grp
SEAGATE ST15150W	/dev/dsk/c0t6d0	8/12.6.0	4095	/dev/vg00
non-LVM disk	Device file	HW Addr	size	swap
SEAGATE ST32430N	/dev/dsk/clt4d0	8/16/5.4.0	2003	0

## DISK AND LVM LAYOUT

## File System layout

LVM Device file	mount point	size	fs type
/dev/vg00:			
/dev/vg00/lvol3	/	84	hfs
/dev/vg00/lvol2	swap	256	
/dev/vg00/lvol1	/stand	48	hfs
/dev/vg00/lvol7	/usr	700	hfs
/dev/vg00/lvol5	/opt	1000	hfs
/dev/vg00/lvol8	/var	160	hfs
/dev/vg00/lvol6	/tmp	32	hfs
/dev/vg00/lvol4	/home	1800	hfs
/dev/vg00	unallocated	12	
Device file	mount point	size	fs type
/dev/dsk/clt4d0	/mnt/clt4d0	2003	hfs

## Swap configuration

type	size	priority	device/location
dev	256	1	/dev/vg00/lvol2

## KERNEL CONFIGURATION

## Kernel Configuration

The following drivers or parameters are configured into your system's kernel. After installing HP-UX, use the sam(lm) command to configure the following items into the kernel:

STRMSGSZ	65535
default_disk_ir	1
maxdsiz	0X10000000
maxssiz	0X10000000
maxtsiz	0X10000000
nstrpty	60

## NETWORK CONFIGURATION

## System Information

The following parameters were set on the configured target:

hostname:	hp46t250
IP address:	15.14.120.250
subnet mask:	255.255.248.0
gateway IP address:	15.14.120.250
time zone:	PST8PDT
DNS domain name:	cup.hp.com
DNS IP address:	15.13.185.120
DNS IP address:	15.13.192.134



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# HP-UX 11.0 Networking

***DAL***

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- Stream s-based TCP/IP
- NFS PV 3
- NIS+
- IP Addressing/Virtual Hosting
- Web Server Support
- Java Support

# Stream s-based TCP/IP

- At 10.30 the TCP/IP stack was changed to stream s-based technology, replacing the BSD-UNIX base
- Very few changes to users or APIs
- Advantages in MP and supportability
- Visible to user as str\* kernel daemons

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# NFS Protocol Version 3

## Advantages

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- Added Support for Large Files
  - Improved Performance
  - Enhanced File Access control
  - New APIs
- 
- Delivered in 10.30 and 11.0 Releases
  - Workstation ACE 2 Release for 10.20 OS

# NFS 3 Support for Large Files

***DAL***

	NFS Version 2	NFS Version 3
File Size and Offsets	32 bits	64 bits
Maximum File Size Supported	2 GBytes	128 GBytes

# NFS PV 3

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## Performance Improvements

- Function calls now return attributes to reduce subsequent `getattr()` function calls
- Read/write blocks can be larger than the previous 8 K byte limit
- Performance of asynchronous write operations much improved: write request from client returns immediately, commit request (NEW) from client causes update to disk on server
- Weak cache consistency: if client modification time matches server modification time the client cache is assumed to be valid
- Remains fully interoperational with NFS V 2

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# N IS+

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- Improved performance and security over N IS
- Continues to support N IS and local file access
- Capability to add future services online without rebooting or reconfiguring
- Delivered in the 10.30 and 11.0 Releases

# NIS+ vs. NIS

## Comparison of Features

***DAL***

	<b>NIS</b>	<b>NIS+</b>
Map updates (Master to Slaves)	Entire map updated	Incremental changes updated
Update propagation	Manual	Automatic
Data Access Restrictions	None	Access controlled on per-entry basis
Authentication	None	Secure RPC
Administration	Must be done on single Master Server	May be made within hierarchical namespace
Namespace	Flat	Hierarchical
Contacting Servers	UDP broadcast	coldstart config file and directory cache

# IP Address Aliasing and Virtual Hosting

***DAL***

- `ifalias` command on 10.20
  - allow multiple IP addresses on same interface  
(IP Address Aliasing)
  - (install appropriate kernel and command patches)
- For 10.30 and 11.0, new options to the `ifconfig` command supply the same functionality
- `ifalias` is obsolete at 10.30 and 11.0



# HP-UX Internet Products

- Netscape FastTrack Server 3.01 bundled with HP-UX
- Netscape SuiteSpot available as separate product:
  - Enterprise Server 3.6 (web server)
  - Messaging Server 3.5 (mail)
  - Directory Server 4.0 (LDAP services)
  - Certificate Server
  - Calendar Server
  - Proxy Server
  - Admin Server 3.5
- Zeus Web Server 3.3
- New HP-UX /I Internet Bundle of products

# Java Support on H P-U X

- Current releases: J D K 1.1.8 and 1.2
  - contains JIT compiler, JVM , class libraries
  - 10.20: uses Java internal (Green) threads
  - 11.0: uses H P-U X kernel threads or Java threads
- JVM 1.2 available in both H otSpot and C lassic versions

# Java Development on HP-UX

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# Java on HP-UX

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- Java Web Page for HP-UX
  - <http://www.hp.com/go/java>
- Can download the Java Development Kit
- Additional documentation available
- Pointers to Java Reference Documents

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# Additional Networking Enhancements

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- New FTP (June 1998 - Patch PHNE\_14479)
  - improved logging, security, on the fly compression
- Kerberos v5 1.0
  - Provides encryption and authentication; simplified installation
- DNS 4.9.6
  - Eases load balancing through round robin
- Sendmail 8.8.6
  - Prevents system overload with anti-spamming feature
- BIND 4.9
  - Improves response times and enhances security
- Gateway Daemon 3.5.1 (gated)
  - Guaranteed connectivity

# HP-UX 11.0

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## Extension Software

Extension Software/Support Plus releases are patch bundles consisting of a set of recommended HP-UX OS patches:

Release Date	OS is Refreshed	Part Number
SP47 December 99	10.20 11.00	B 3782-10453
SP46 September 99	10.20 11.00	B 3782-10438
XR 45 June 99	10.10 10.20 11.00	B 3782-10418
XR 44 April 99	10.10 10.20 11.00	B 3782-10386
XR 43 February 99	10.10 10.20 11.00	B 3782-10371
XR 41 October 98	10.10 10.20 11.00	B 3782-10356
XR 40 August 98	10.10 10.20 11.00	B 3782-10340
XR 39 June 98	10.10 10.20 11.00	B 3782-10346
XR 38 April 98	10.10 10.20 11.00	B 3782-10300
XR 37 February 98	10.01 10.10 11.00	B 3782-10293

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# HP-UX 11

## Instant Ignition Bundle

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### Core Products:

- Netscape Fasttrack Server
- Web QoS Peak Service
- JAVA (JRE, JDK, currently instantly ignitable)
- CD SA Framework
- Network Drivers (all currently instantly ignitable)
- Ignite/UX
- EMS
- Kernel optimizations for web performance

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# HP-UX 11 Product Bundle

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## Optional Products:

- Manageability                      OpenView, Service Control  
    WebQoS Premium
- Web Servers                              Netscape, Zeus, Apache
- Availability                              MC/ServiceGuard
- Security                                      Praesidium
- App. Dev.                                      BEA
- Services                                      Pre and Post sales